<u>scal.cs.uni-</u> <u>salzburg.at</u>

multicore-scalable concurrent data

structures

<u>scalloc.cs.uni-</u> <u>salzburg.at</u>

multicore-scalable concurrent allocator

<u>selfie.cs.uni-</u>

salzburg.at

self-referential systems software for teaching

Scal, Scalloc, and Selfie

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UPMARC Multicore Computing Summer School, Uppsala University, June 2016

Joint Work

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Concurrent Data Structures: <u>scal.cs.uni-salzburg.at</u> [POPL13, CF13, POPL15, NETYS15]

- Scal is an open-source benchmarking framework that provides
 - 1. software <u>infrastructure</u> for executing concurrent data structure algorithms,
 - 2. workloads for <u>benchmarking</u> their performance and scalability, and
 - 3. <u>implementations</u> of a large set of concurrent data structures.

Scal: A Benchmarking Suite for Concurrent Data Structures [NETYS15]

Name	Semantics	Year	Ref
Lock-based Singly-linked	strict queue	1968	[1]
Michael Scott (MS) Queue	strict queue	1996	[2]
Flat Combining Queue	strict queue	2010	[3]
Wait-free Queue	strict queue	2012	[4]
Linked Cyclic Ring Queue	strict queue	2013	[5]
Timestamped (TS) Queue	strict queue	2015	[6]
Cooperative TS Queue	strict queue	2015	[7]
Segment Queue	k-relaxed queue	2010	[8]
Random Dequeue (RD)	k-relaxed queue	2010	[8]
Bounded Size k-FIFO	k-relaxed queue, pool	2013	[9]
Unbounded Size k-FIFO	k-relaxed queue, pool	2013	[9]
b-RR Distributed Queue	k-relaxed queue, pool	2013	[10]
Least-Recently-Used (LRU)	k-relaxed queue, pool	2013	[10]
Locally Linearizable DQ	locally linearizable	2015	[11]
Locally Linearizable k-FIFO	locally linearizable	2015	[11]
Relaxed TS Queue	quiescently consistent	2015	[7]
Lock-based Singly-linked	strict stack	1968	[1]
Treiber Stack	strict stack	1986	[12]
Elimination-backoff Stack	strict stack	2004	[13]
Timestamped (TS) Stack	strict stack	2015	[6]
k-Stack	k-relaxed stack	2013	[14]
b-RR Distributed Stack (DS)	k-relaxed stack, pool	2013	[10]
Least-Recently-Used (LRU)	k-relaxed stack, pool	2013	[10]
Locally Linearizable DS	locally linearizable	2015	[11]
Locally Linearizable k-Stack	locally linearizable	2015	[11]
Timestamped (TS) Deque	strict deque	2015	[7]
d-RA DQ and DS	strict pool	2013	[10]



k-FIFO Queues [PaCT13]



(a) Very high contention (c = 1000, i = 0)





Distributed Queues [CF13]



(a) High contention producer-consumer microbenchmark (c = 250) (b) Low contention producer-consumer microbenchmark (c = 2000) Figure 1: Performance and scalability of producer-consumer microbenchmarks with an increasing number of threads on a 40-core (2 hyper-threads per core) server machine

Timestamped (TS) Stack [POPL15]



(a) Producer-consumer benchmark, 40-core machine.

(b) Producer-consumer benchmark, 64-core machine.

Local Linearizability [CONCUR16]



"queue-like" data structures

"stack-like" data structures

Figure 5 Performance and scalability of producer-consumer microbenchmarks with an increasing number of threads on a 40-core (2 hyperthreads per core) machine

Scalloc: Concurrent Memory Allocator scalloc.cs.uni-salzburg.at [OOPSLA15]

- fast, multicore-scalable, low-memory-overhead allocator
- three key ideas:
 - 1. backend: <u>single</u> global concurrent data structure for reclaiming memory <u>effectively</u> and <u>efficiently</u>
 - 2. virtual spans: <u>single</u> algorithm for <u>small</u> and <u>big</u> objects
 - 3. frontend: constant-time (modulo synchronization) allocation and <u>eager</u> deallocation

Local Allocation & Deallocation

Figure 6: Thread-local workload: Threadtest benchmark

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Remote Deallocation

Figure 9: Temporal and spatial performance for the producer-consumer experiment

Figure 10: Temporal and spatial performance for the object-size robustness experiment at 40 threads

Memory Access

Figure 11: Memory access time for the locality experiment

Virtual Spans: 64-bit Address Space

Figure 1: Structure of arena, virtual spans, and real spans

Backend: Double Segregation

Figure 2: Span pool layout

Frontend: Eager Memory Reuse

Figure 3: Life cycle of a span

Selfie: Teaching Systems Engineering [selfie.cs.uni-salzburg.at]

- * *Selfie* is a self-referential 6k-line C implementation (in a single file) of:
 - 1. a <u>self-compiling</u> compiler called *starc* that compiles a tiny subset of C called C Star (C*) to a tiny subset of MIPS32 called MIPSter,
 - 2. a <u>self-executing</u> emulator called *mipster* that executes MIPSter code including itself when compiled with starc,
 - 3. a <u>self-hosting</u> hypervisor called *hypster* that virtualizes mipster and can host all of selfie including itself, and
 - 4. a tiny C* library called *libcstar* utilized by starc and mipster.

5 statements: assignment while if return procedure()

return n;

library: exit, malloc, open, read, write

MIPSter: 17 out of 43 Instructions

atoi.c: \$pc=0x000001CC: lw \$t0,-4(\$fp) \$pc=0x000001D0: addiu \$t1,\$zero,1 atoi.c: i = i + 1;atoi.c: \$pc=0x000001D4: addu \$t0,\$t0,\$t1 \$pc=0x000001D8: sw \$t0,-4(\$fp) atoi.c: \$pc=0x000001DC: lw \$t0,8(\$fp) atoi.c: \$pc=0x000001E0: lw \$t1,-4(\$fp) atoi.c: \$pc=0x000001E4: addiu \$t2,\$zero,4 atoi.c: \$pc=0x000001E8: multu \$t1,\$t2 atoi.c: \$pc=0x000001EC: mflo \$t1 atoi.c: c = *(s + i);atoi.c: \$pc=0x000001F0: nop atoi.c: \$pc=0x000001F4: nop atoi.c: \$pc=0x000001F8: addu \$t0,\$t0,\$t1 lw \$t0,0(\$t0) atoi.c: \$pc=0x000001FC: sw \$t0,-12(\$fp) \$pc=0x00000200: atoi.c:

Future Work with Selfie et al.

- ✤ I/O
- file systems
- memory allocation
- garbage collection
- concurrency: <u>semantics</u>
- parallelism: <u>multicore</u>
- volatility: <u>persistent</u> memory

