

scal.cs.uni-
salzburg.at

multicore-scalable
concurrent data
structures

scaloc.cs.uni-
salzburg.at

multicore-scalable
concurrent allocator

selfie.cs.uni-
salzburg.at

self-referential systems
software for teaching

Scal, Scalloc, and Selfie

Christoph Kirsch, University of Salzburg, Austria

Joint Work

- ❖ Martin Aigner
- ❖ Christian Barthel
- ❖ Mike Dodds
- ❖ Andreas Haas
- ❖ Thomas Henzinger
- ❖ Andreas Holzer
- ❖ Thomas Hütter
- ❖ Michael Lippautz
- ❖ Alexander Miller
- ❖ Simone Oblasser
- ❖ Hannes Payer
- ❖ Mario Preishuber
- ❖ Ana Sokolova
- ❖ Ali Szegin

Concurrent Data Structures: scal.cs.uni-salzburg.at [POPL13, CF13, POPL15, NETYS15]

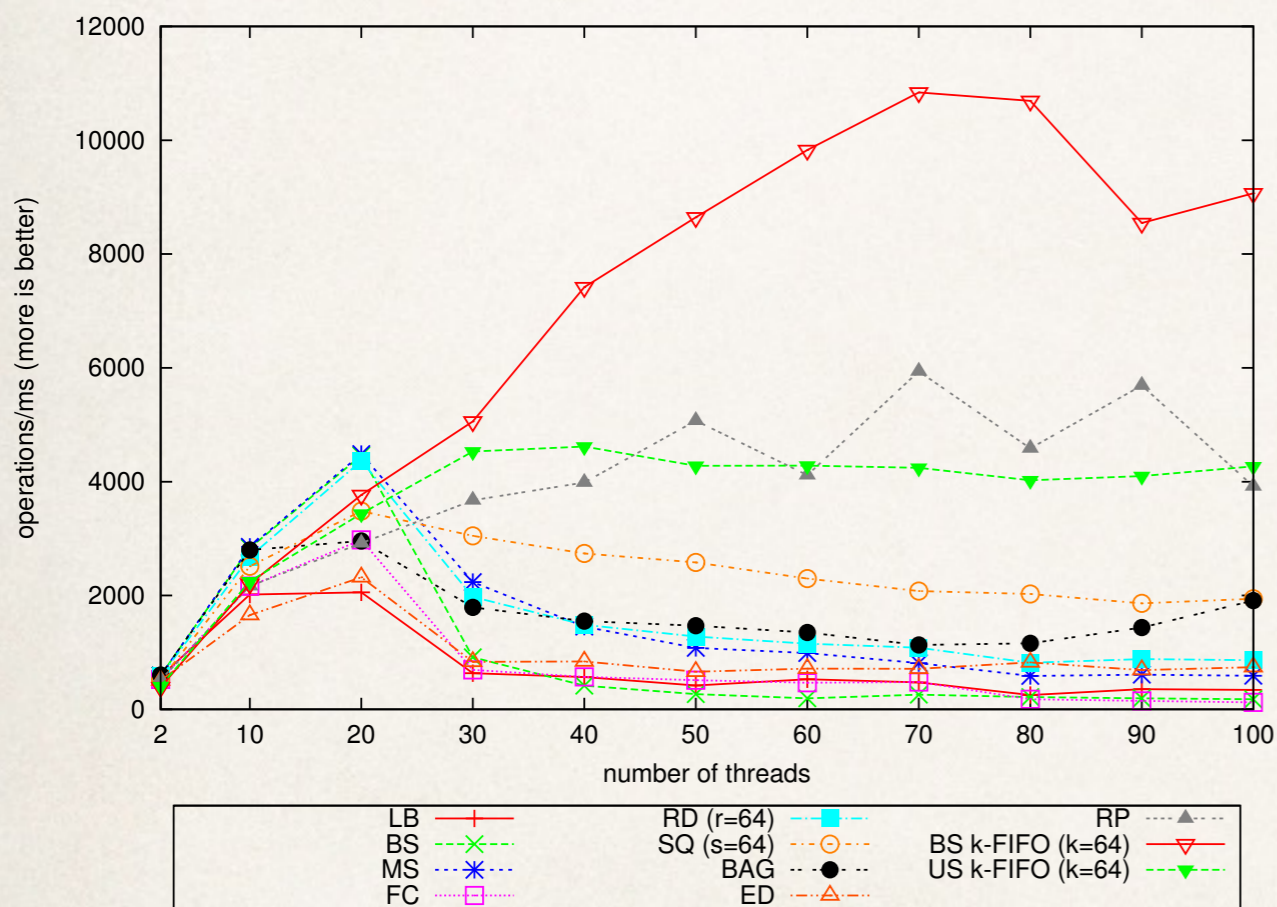
- ❖ Scal is an open-source benchmarking framework that provides
 1. software infrastructure for executing concurrent data structure algorithms,
 2. workloads for benchmarking their performance and scalability, and
 3. implementations of a large set of concurrent data structures.

Scal: A Benchmarking Suite for Concurrent Data Structures [NETYS15]

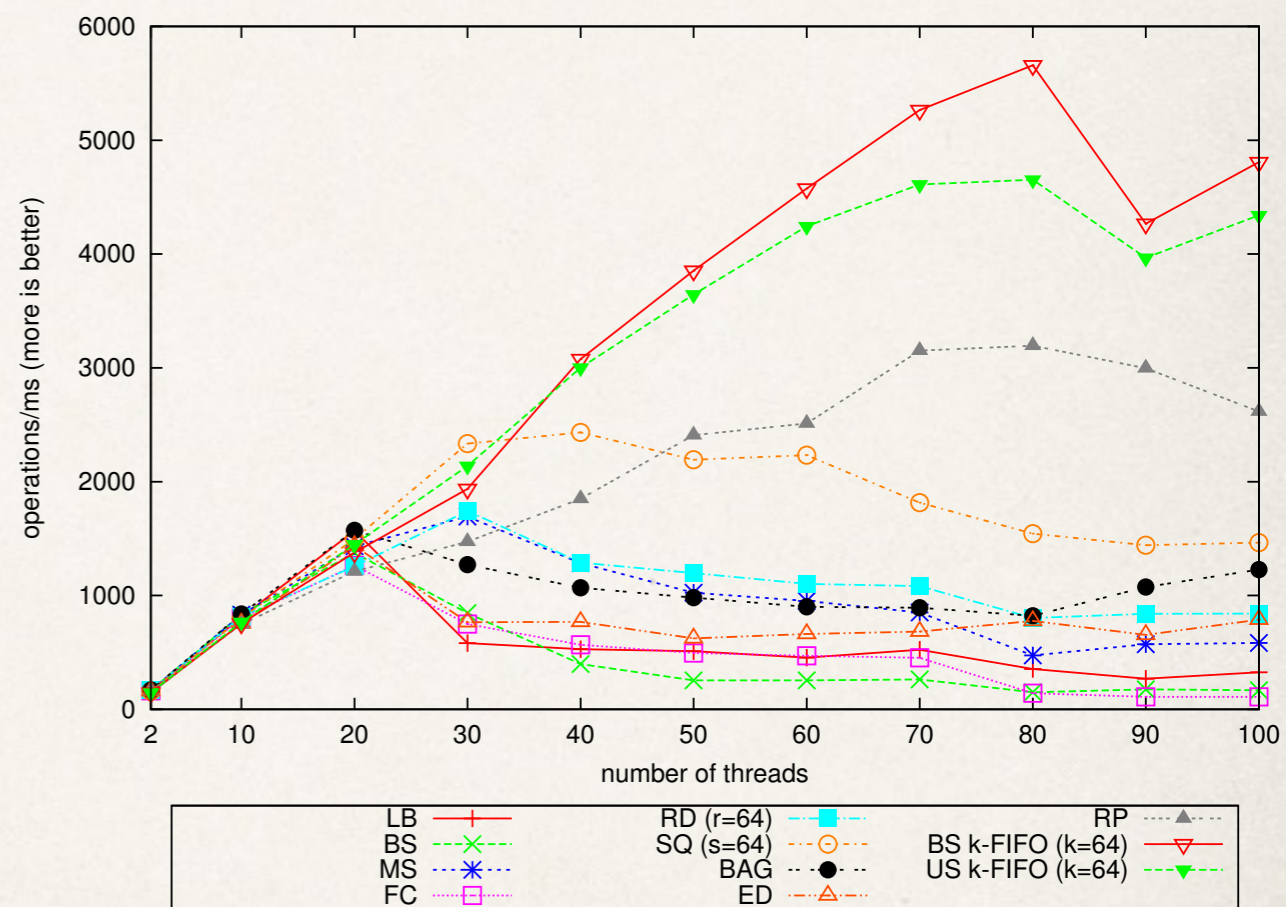
Name	Semantics	Year	Ref
Lock-based Singly-linked	strict queue	1968	[1]
Michael Scott (MS) Queue	strict queue	1996	[2]
Flat Combining Queue	strict queue	2010	[3]
Wait-free Queue	strict queue	2012	[4]
Linked Cyclic Ring Queue	strict queue	2013	[5]
Timestamped (TS) Queue	strict queue	2015	[6]
Cooperative TS Queue	strict queue	2015	[7]
Segment Queue	k-relaxed queue	2010	[8]
Random Dequeue (RD)	k-relaxed queue	2010	[8]
Bounded Size k-FIFO	k-relaxed queue, pool	2013	[9]
Unbounded Size k-FIFO	k-relaxed queue, pool	2013	[9]
b-RR Distributed Queue	k-relaxed queue, pool	2013	[10]
Least-Recently-Used (LRU)	k-relaxed queue, pool	2013	[10]
Locally Linearizable DQ	locally linearizable	2015	[11]
Locally Linearizable k-FIFO	locally linearizable	2015	[11]
Relaxed TS Queue	quiescently consistent	2015	[7]
Lock-based Singly-linked	strict stack	1968	[1]
Treiber Stack	strict stack	1986	[12]
Elimination-backoff Stack	strict stack	2004	[13]
Timestamped (TS) Stack	strict stack	2015	[6]
k-Stack	k-relaxed stack	2013	[14]
b-RR Distributed Stack (DS)	k-relaxed stack, pool	2013	[10]
Least-Recently-Used (LRU)	k-relaxed stack, pool	2013	[10]
Locally Linearizable DS	locally linearizable	2015	[11]
Locally Linearizable k-Stack	locally linearizable	2015	[11]
Timestamped (TS) Deque	strict deque	2015	[7]
d-RA DQ and DS	strict pool	2013	[10]



k-FIFO Queues [PaCT13]

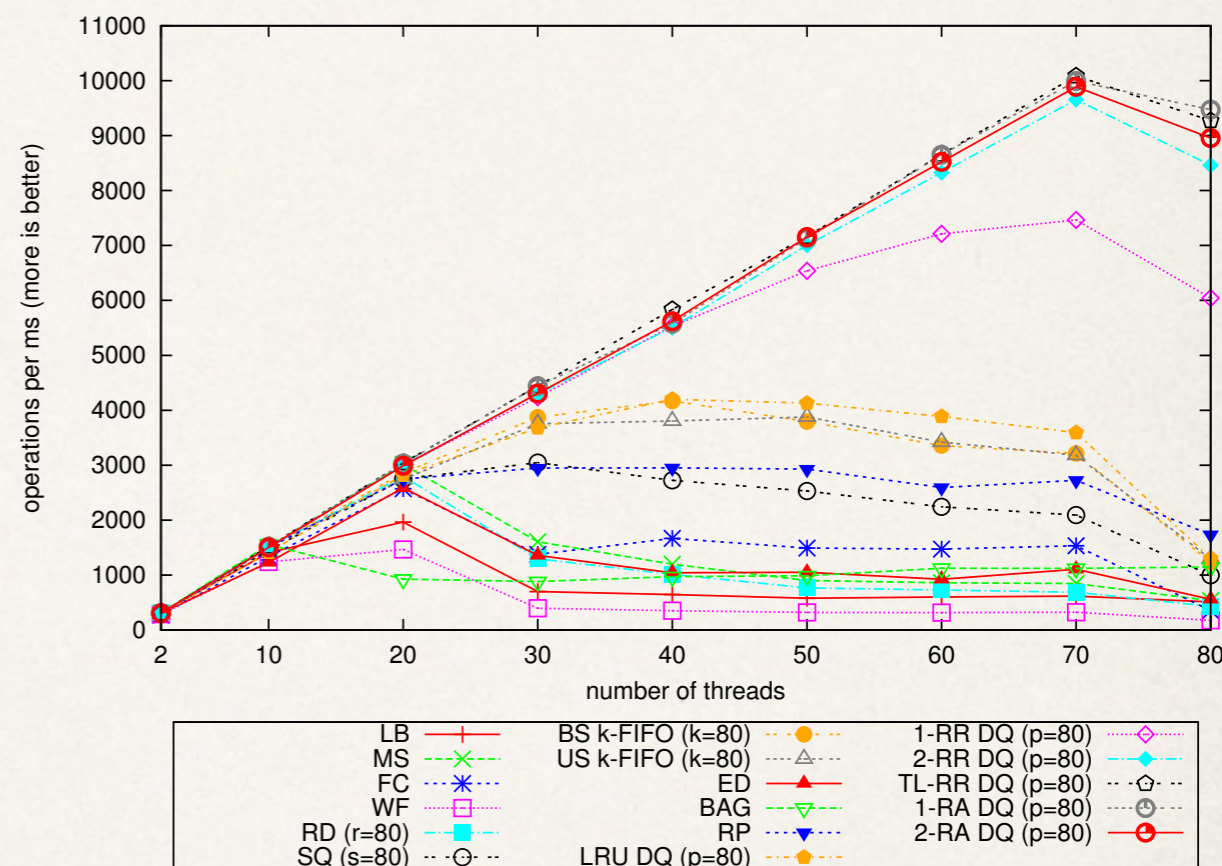
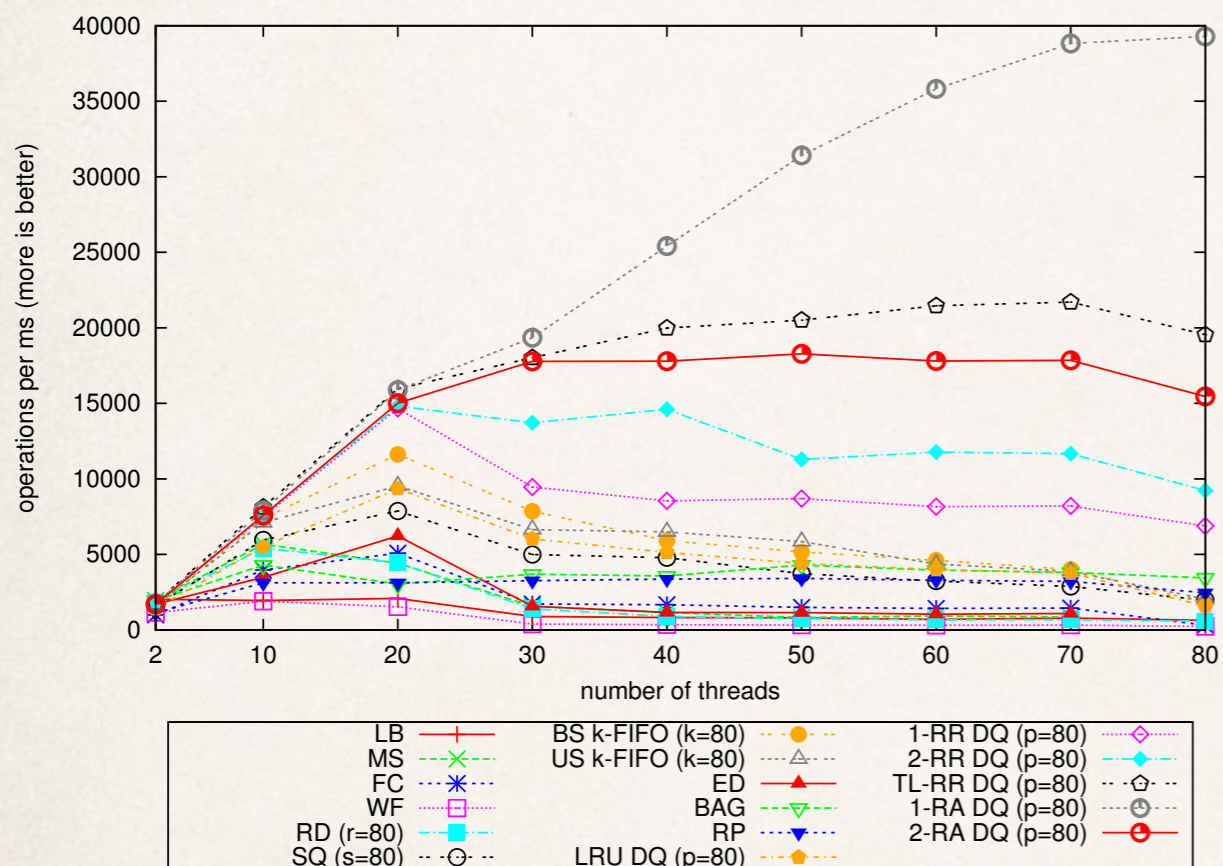


(a) Very high contention ($c = 1000, i = 0$)



(b) High contention ($c = 4000, i = 0$)

Distributed Queues [CF13]

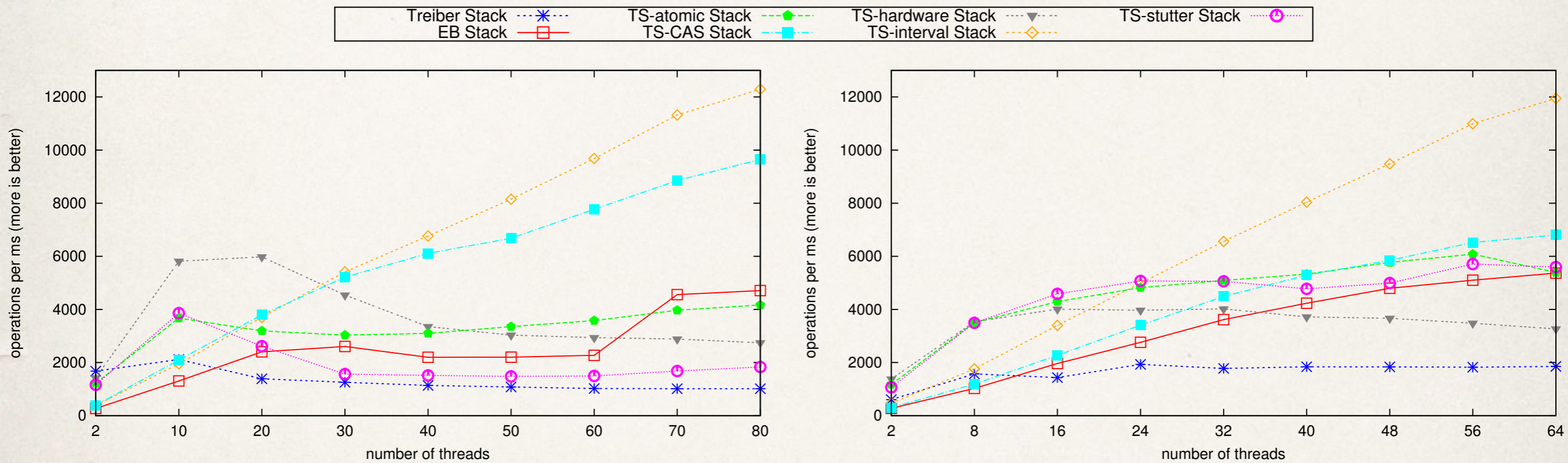


(a) High contention producer-consumer microbenchmark ($c = 250$)

(b) Low contention producer-consumer microbenchmark ($c = 2000$)

Figure 1: Performance and scalability of producer-consumer microbenchmarks with an increasing number of threads on a 40-core (2 hyper-threads per core) server machine

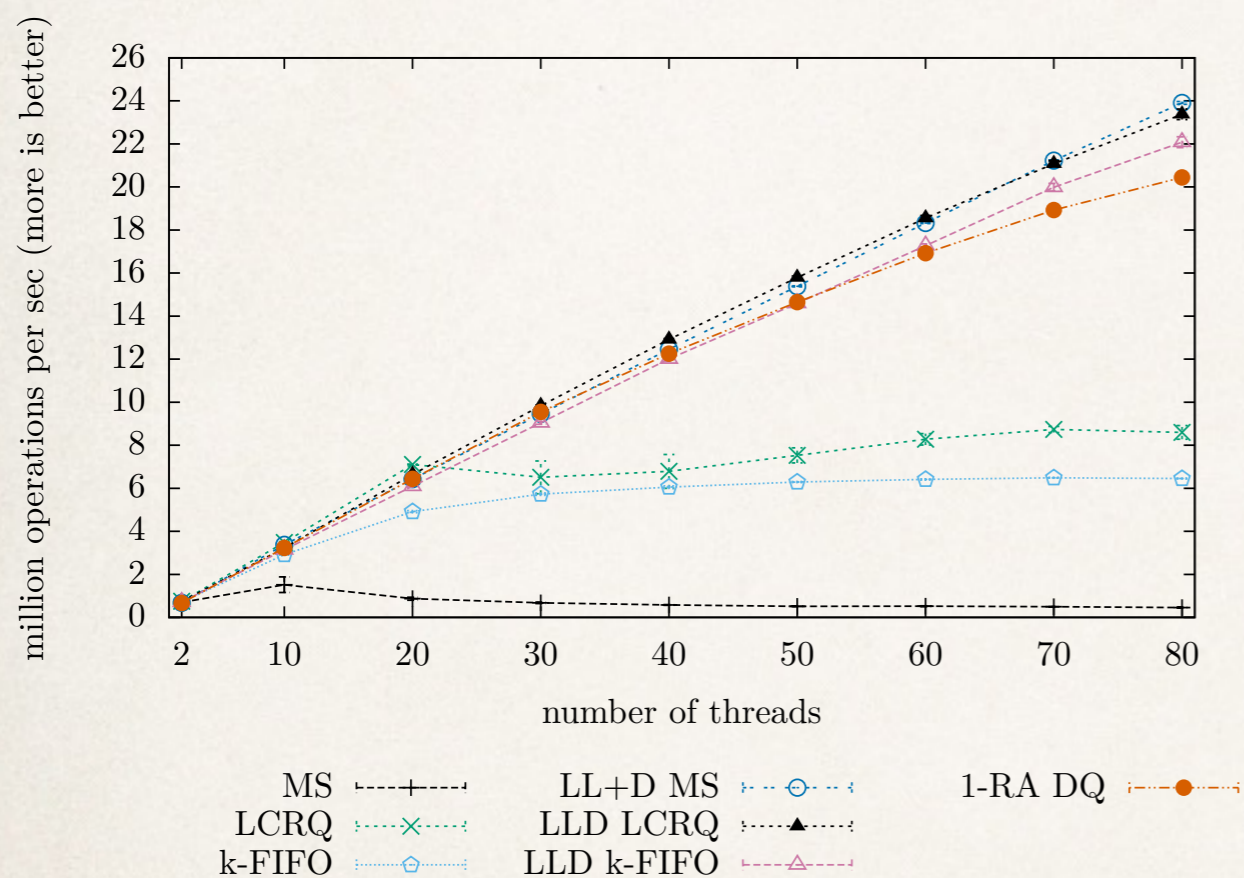
Timestamped (TS) Stack [POPL15]



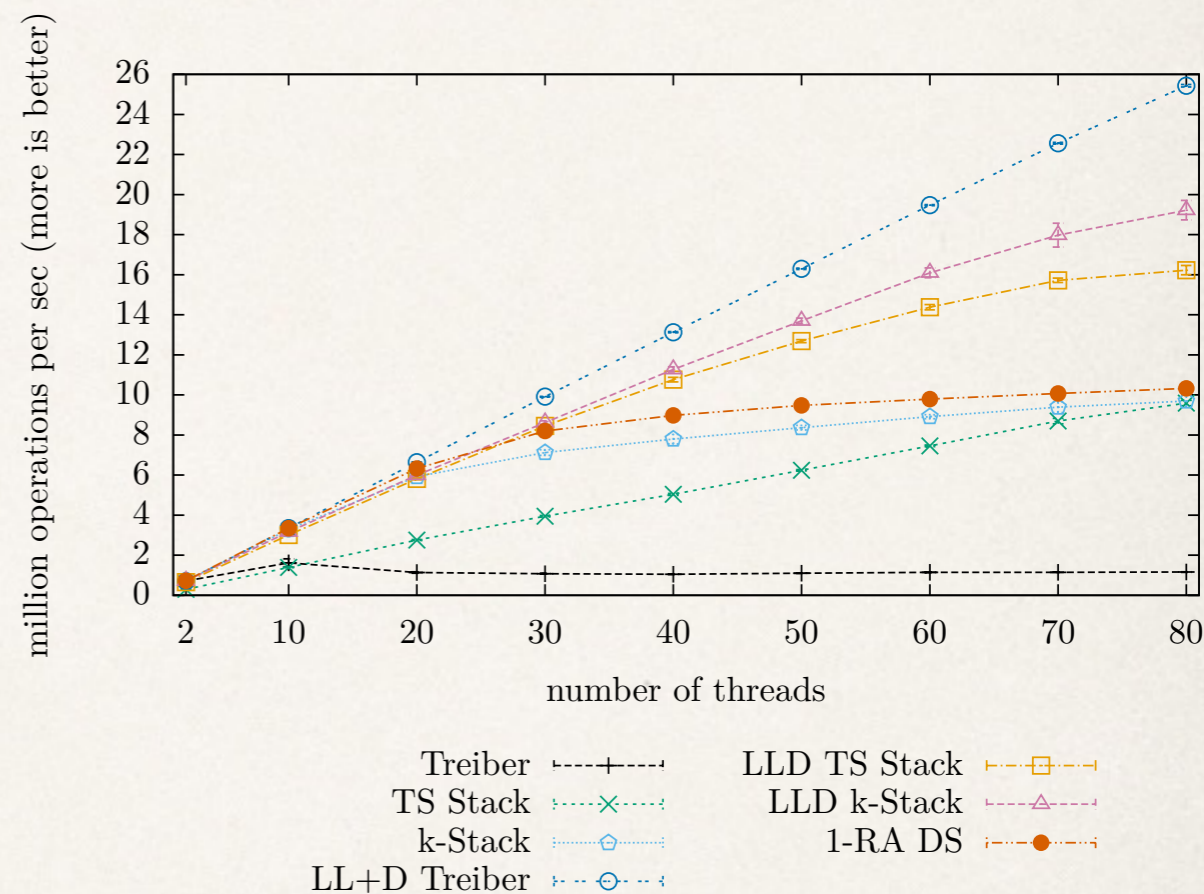
(a) Producer-consumer benchmark, 40-core machine.

(b) Producer-consumer benchmark, 64-core machine.

Local Linearizability [CONCUR16]



“queue-like” data structures



“stack-like” data structures

Figure 5 Performance and scalability of producer-consumer microbenchmarks with an increasing number of threads on a 40-core (2 hyperthreads per core) machine

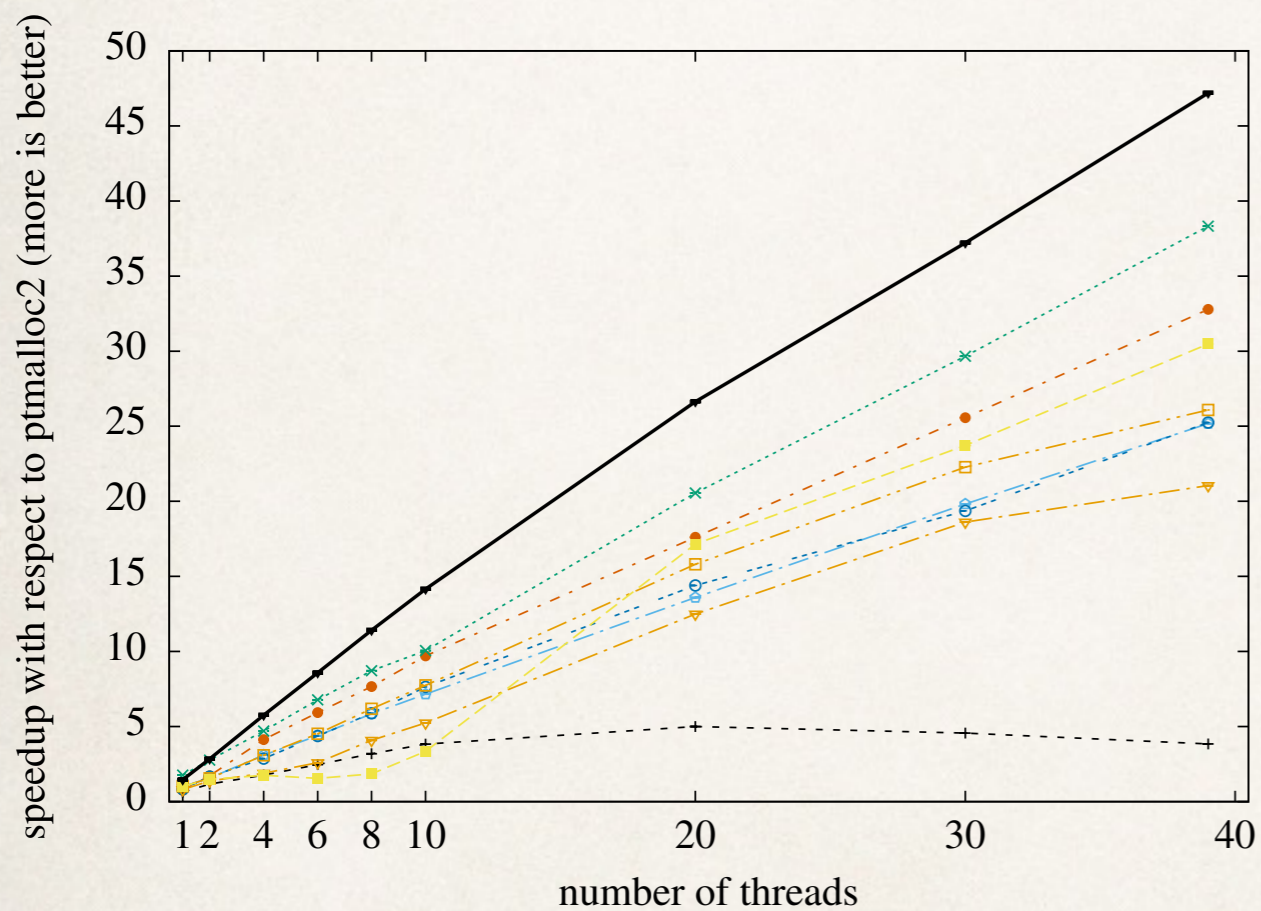


Scalloc: Concurrent Memory Allocator

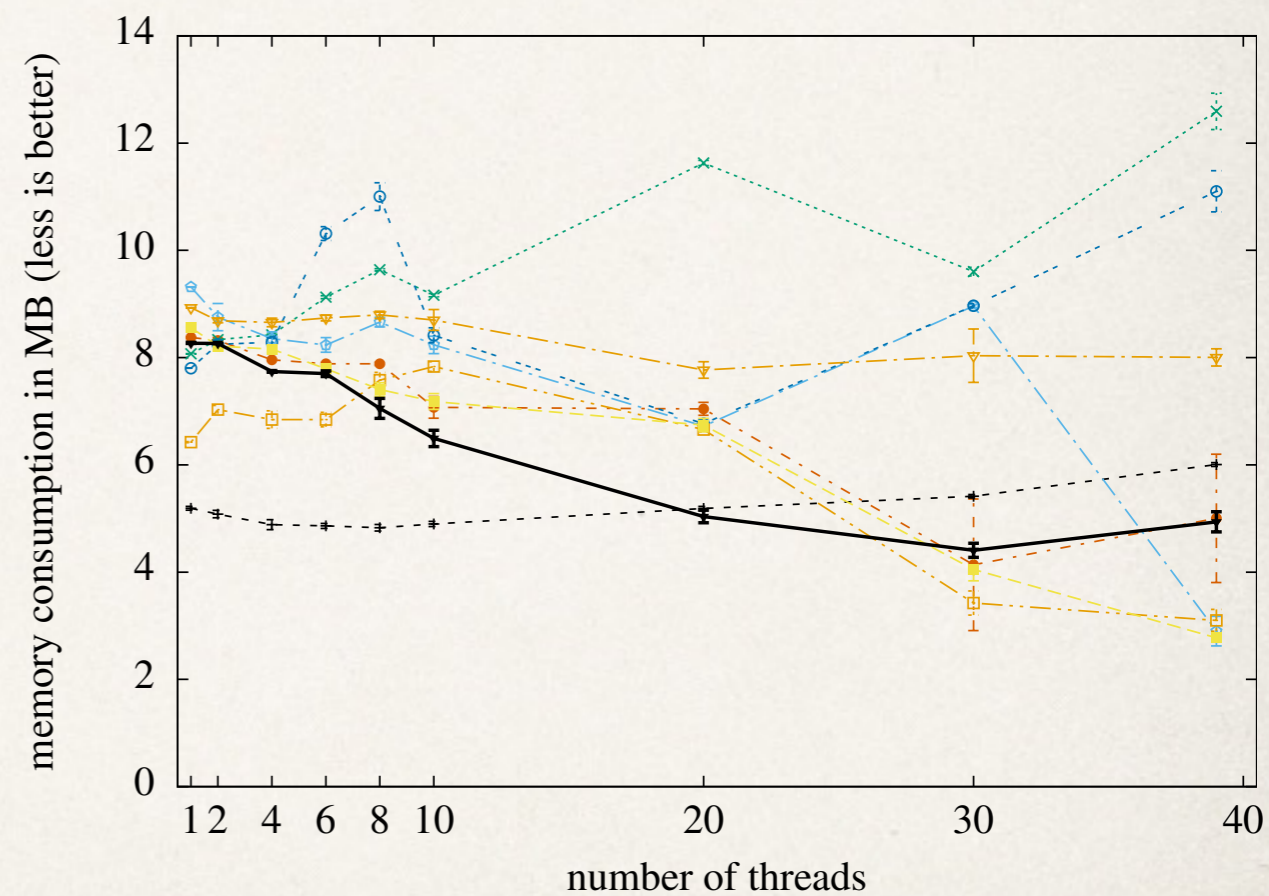
scalloc.cs.uni-salzburg.at [OOPSLA15]

- ❖ fast, multicore-scalable, low-memory-overhead allocator
- ❖ three key ideas:
 1. backend: single global concurrent data structure for reclaiming memory effectively and efficiently
 2. virtual spans: single algorithm for small and big objects
 3. frontend: constant-time (modulo synchronization) allocation and eager deallocation

Local Allocation & Deallocation



(a) Speedup



(b) Memory consumption

Figure 6: Thread-local workload: Threadtest benchmark

Remote Deallocation

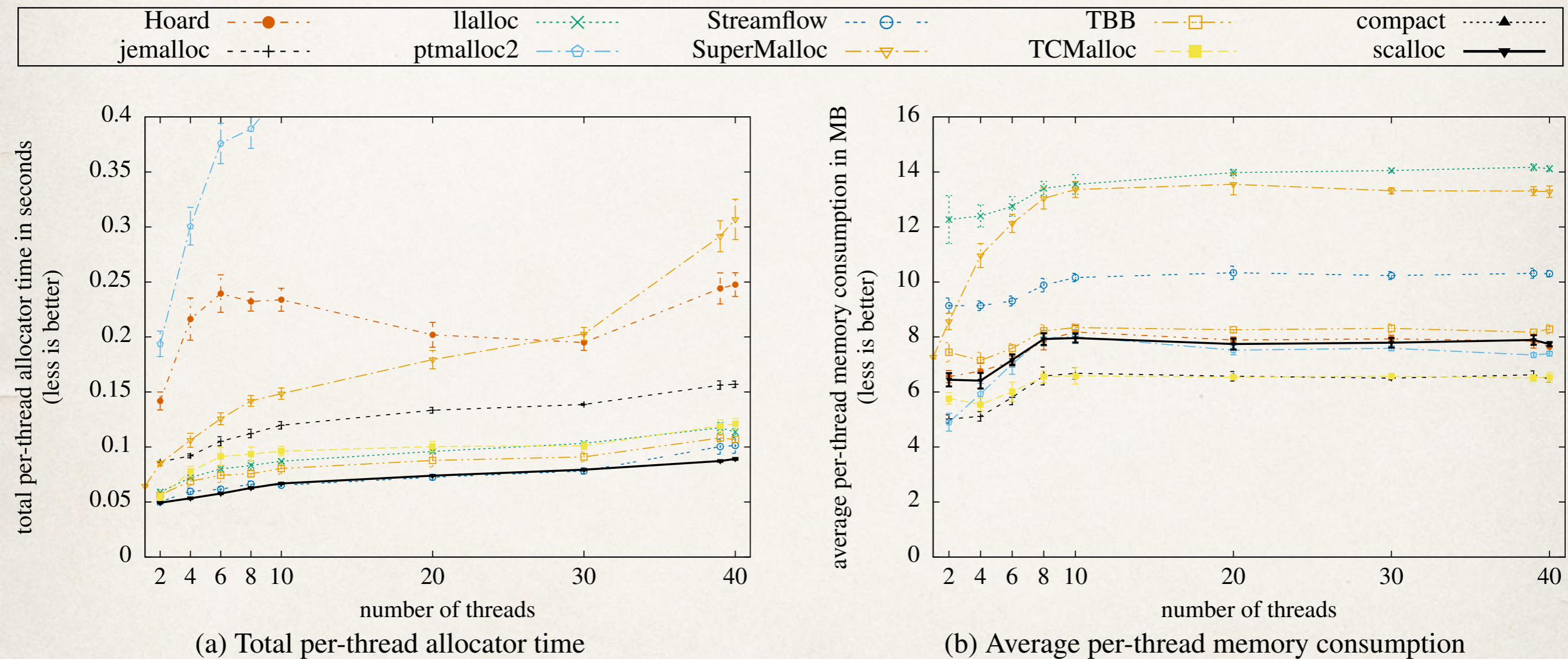


Figure 9: Temporal and spatial performance for the producer-consumer experiment

Object Size

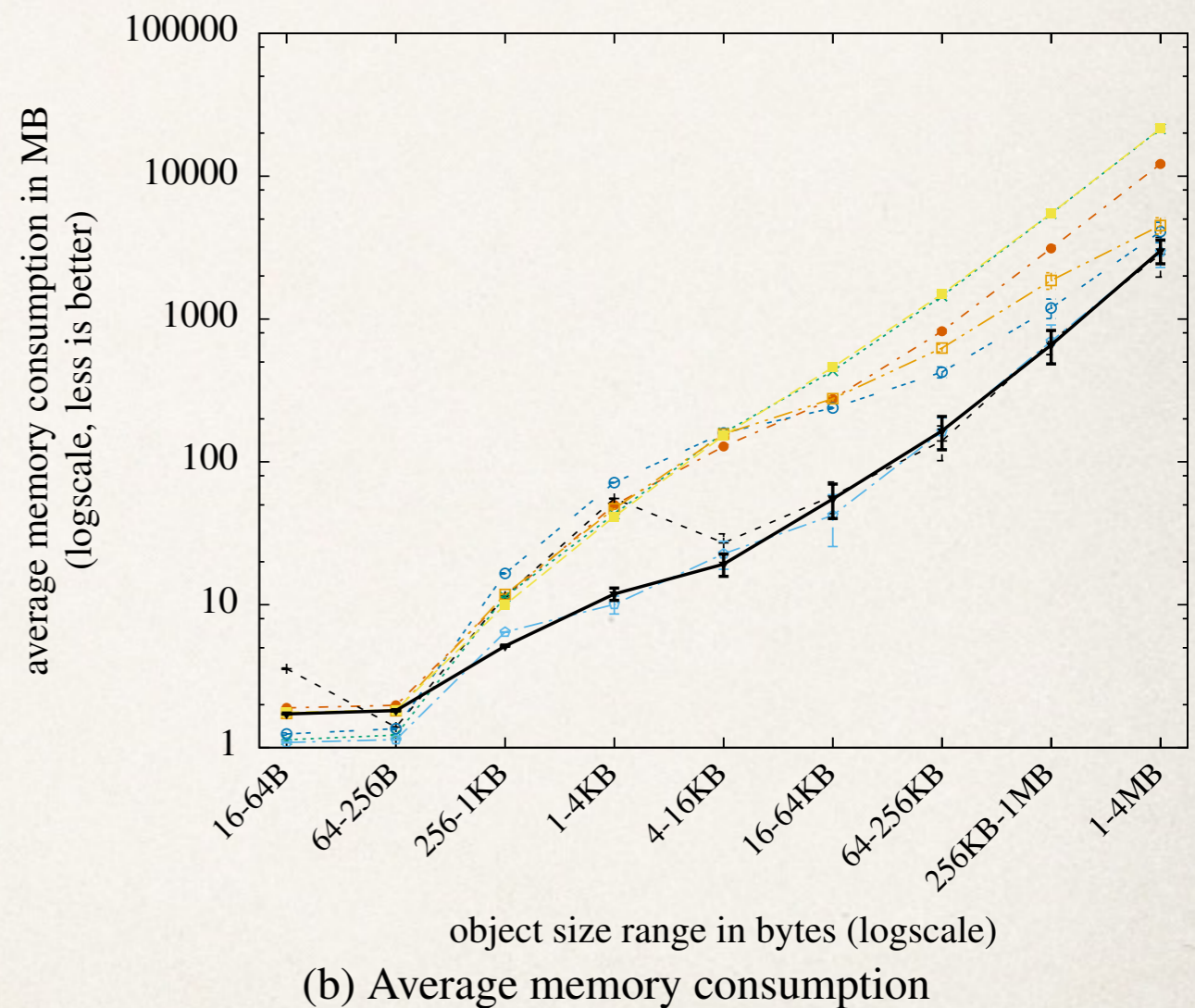
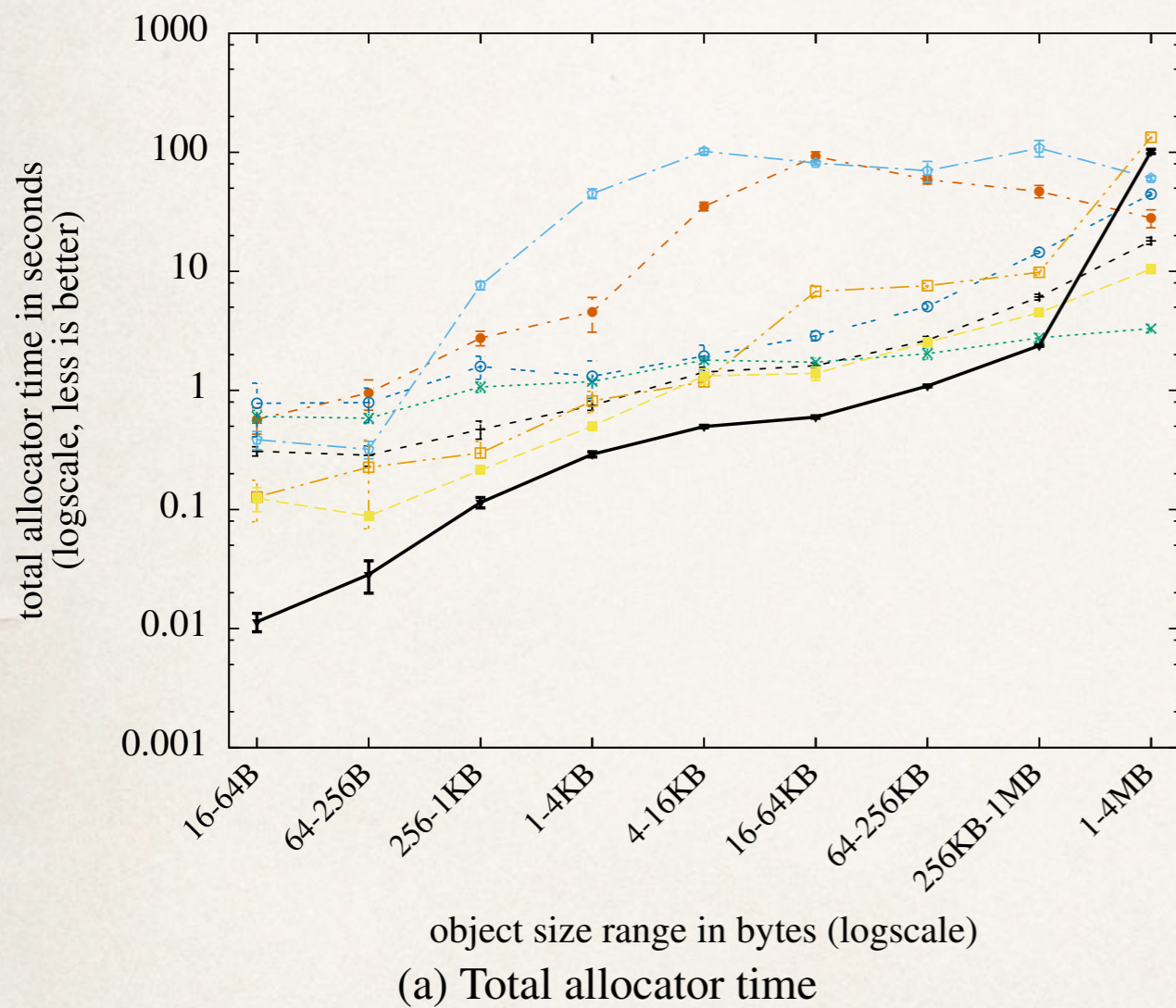


Figure 10: Temporal and spatial performance for the object-size robustness experiment at 40 threads

Memory Access

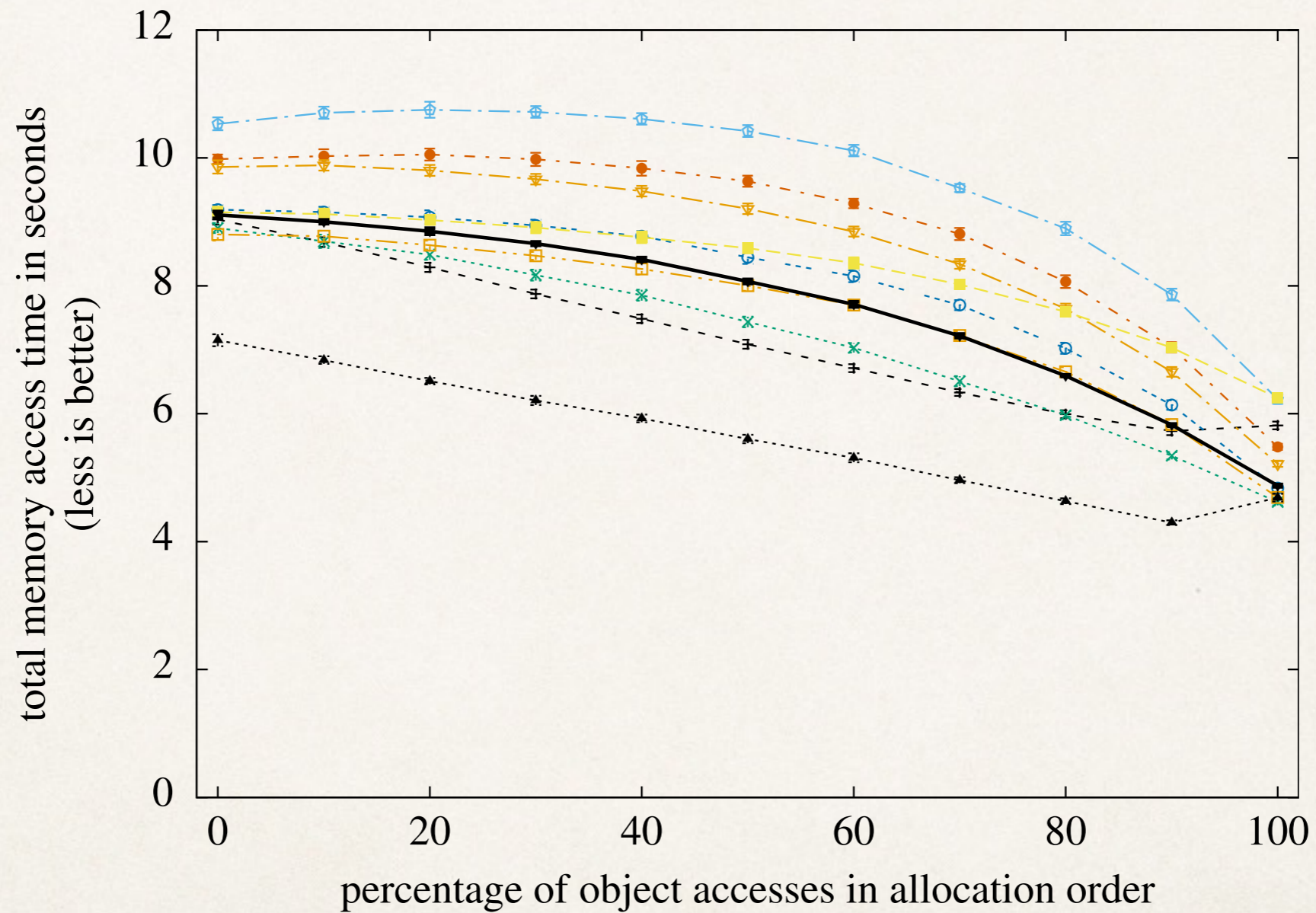


Figure 11: Memory access time for the locality experiment

Virtual Spans: 64-bit Address Space

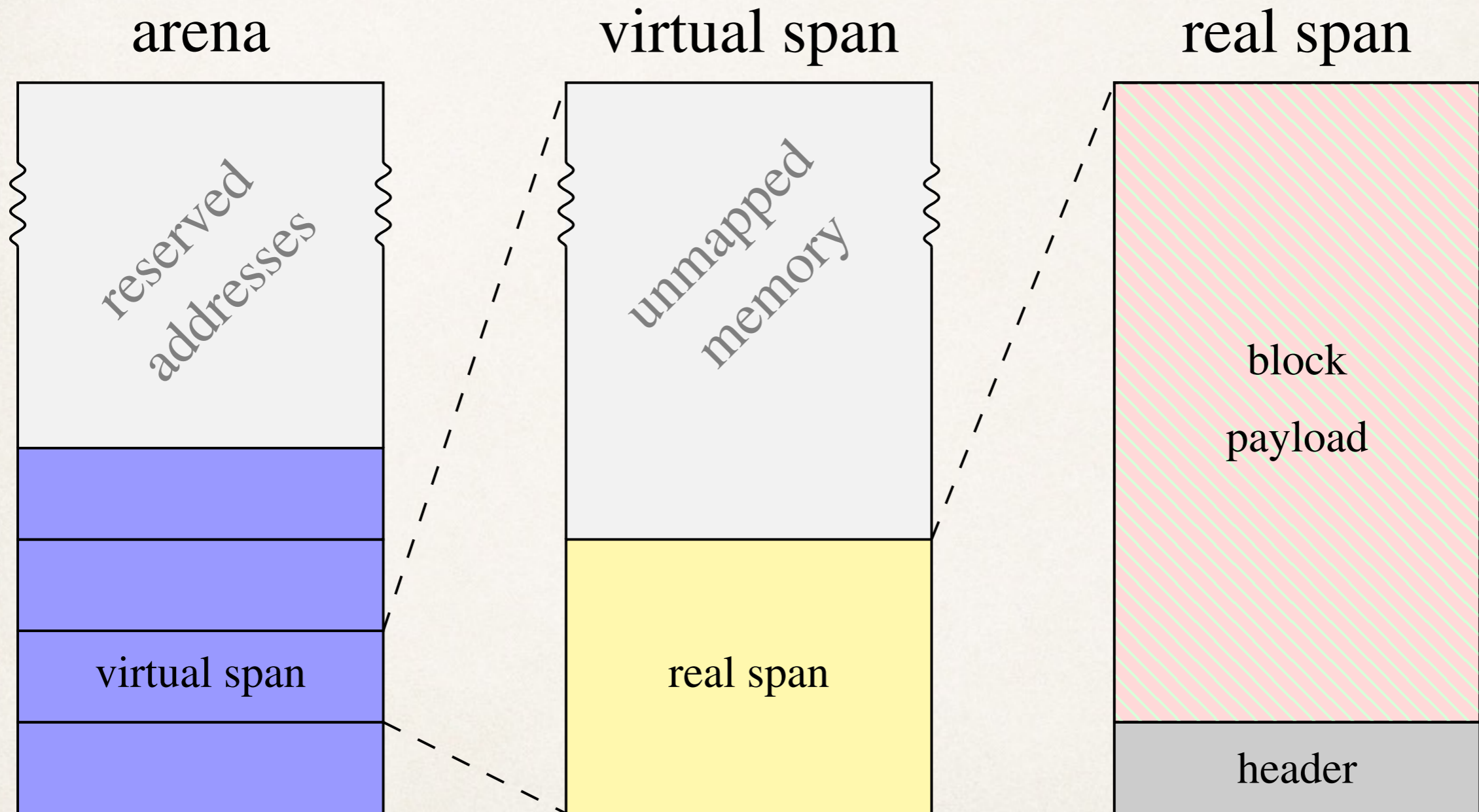


Figure 1: Structure of arena, virtual spans, and real spans

Backend: Double Segregation

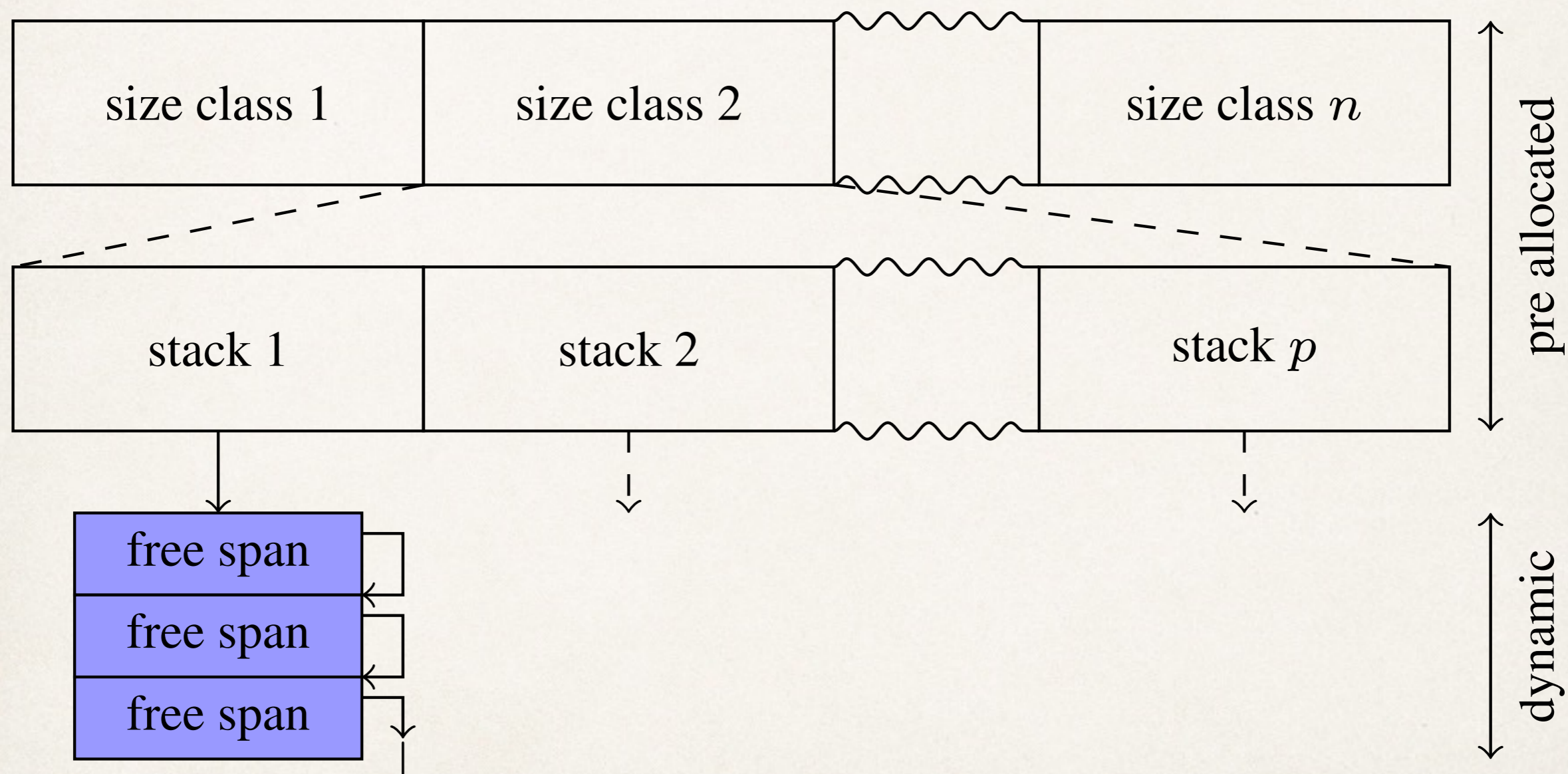


Figure 2: Span pool layout

Frontend: Eager Memory Reuse

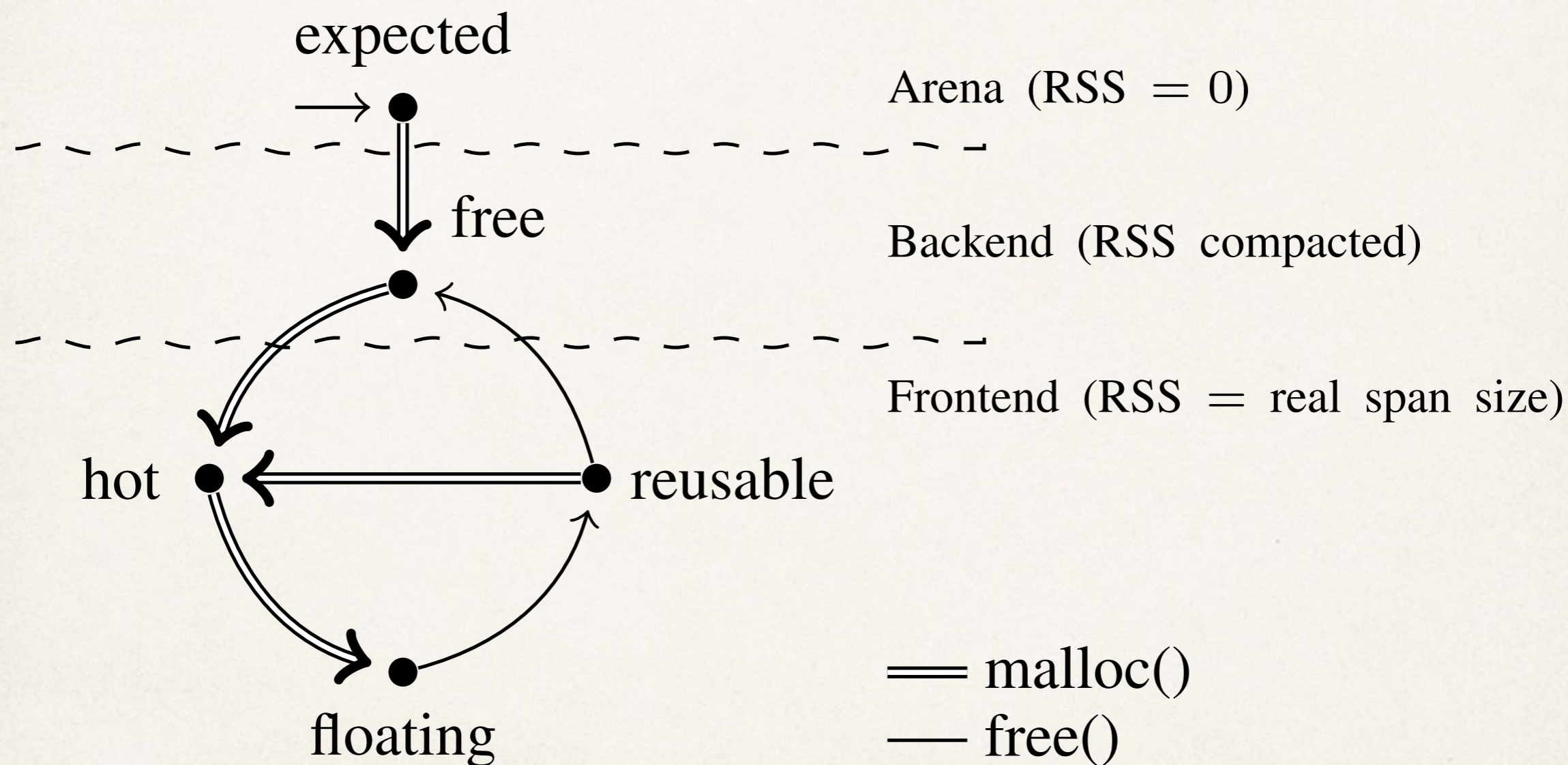


Figure 3: Life cycle of a span







Selfie: Teaching Systems Engineering

[\[selfie.cs.uni-salzburg.at\]](http://selfie.cs.uni-salzburg.at)

- ❖ *Selfie* is a self-referential 6k-line C implementation (in a single file) of:
 1. a self-compiling compiler called *starc* that compiles a tiny subset of C called C Star (C*) to a tiny subset of MIPS32 called MIPSter,
 2. a self-executing emulator called *mipster* that executes MIPSter code including itself when compiled with *starc*,
 3. a self-hosting hypervisor called *hypster* that virtualizes *mipster* and can host all of *selfie* including itself, and
 4. a tiny C* library called *libcstar* utilized by *starc* and *mipster*.

5 statements:
assignment
while
if
return
procedure()

```
int atoi(int *s) {  
    int i;  
    int n;  
    int c;  
  
    i = 0;  
    n = 0;  
    c = *(s+i);
```

no data structures,
just int and int*
and dereferencing:
the * operator

character literals
string literals

```
while (c != 0) {  
    n = n * 10 + c - '0';  
    if (n < 0)  
        return -1;
```

integer arithmetics
pointer arithmetics

```
    i = i + 1;  
    c = *(s+i);
```

no bitwise operators
no Boolean operators

```
return n;
```

library: exit, malloc, open, read, write

MIPSter: 17 out of 43 Instructions

```
atoi.c: $pc=0x000001CC: lw $t0,-4($fp)
atoi.c: $pc=0x000001D0: addiu $t1,$zero,1
atoi.c: $pc=0x000001D4: addu $t0,$t0,$t1
atoi.c: $pc=0x000001D8: sw $t0,-4($fp)
atoi.c: $pc=0x000001DC: lw $t0,8($fp)
atoi.c: $pc=0x000001E0: lw $t1,-4($fp)
atoi.c: $pc=0x000001E4: addiu $t2,$zero,4
atoi.c: $pc=0x000001E8: multu $t1,$t2
atoi.c: $pc=0x000001EC: mflo $t1
atoi.c: $pc=0x000001F0: nop
atoi.c: $pc=0x000001F4: nop
atoi.c: $pc=0x000001F8: addu $t0,$t0,$t1
atoi.c: $pc=0x000001FC: lw $t0,0($t0)
atoi.c: $pc=0x00000200: sw $t0,-12($fp)
```



`i = i + 1;`

`c = *(s + i);`

Future Work with Selfie et al.

- ❖ I/O
- ❖ file systems
- ❖ memory allocation
- ❖ garbage collection
- ❖ concurrency: semantics
- ❖ parallelism: multicore
- ❖ volatility: persistent memory





Thank you!