

Tiptoe: A Compositional Real-Time Operating System

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tiptoe.cs.uni-salzburg.at

- Silviu Craciunas* (Benchmarking)
- Hannes Payer (Memory Management)
- Ana Sokolova* (Theoretical Foundation)
- Horst Stadler (I/O Subsystem)
- Robert Staudinger* (Kernel)

Process A



Process B

Operating System

Memory

CPU

I/O

“Theorem”

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- (Compositionality) The **time** and **space** a software process needs to execute is determined by the **process**, not the system and not other software processes.

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- (Predictability) The **system** can tell how much **time** and **space** is available without looking at any existing software processes.

“Corollary”

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- (Memory) The time a software process takes to **allocate** and **free** a memory object is determined by the size of the **object**.

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- **(Memory)** The time a software process takes to **allocate** and **free** a memory object is determined by the size of the **object**.
- **(I/O)** The time a software process takes to **read** input data and **write** output data is determined by the size of the **data**.

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- 1ms/100ms CPU time (\neq 10ms/s)
- 4MB/2s memory allocation rate
- 1KB/10ms network bandwidth
- 10J/100ms energy consumption

Outline

1. Memory Management
2. Concurrency Management
3. I/O Management

Toe A



Toe B

Tip

Memory

CPU

I/O

Outline

1. Memory Management
2. Concurrency Management
3. I/O Management

Tiptoe System

P2P Ethernet
Connection

OR

Serial
Connection

I/O Host Computer

Network

Disk

AD/DA

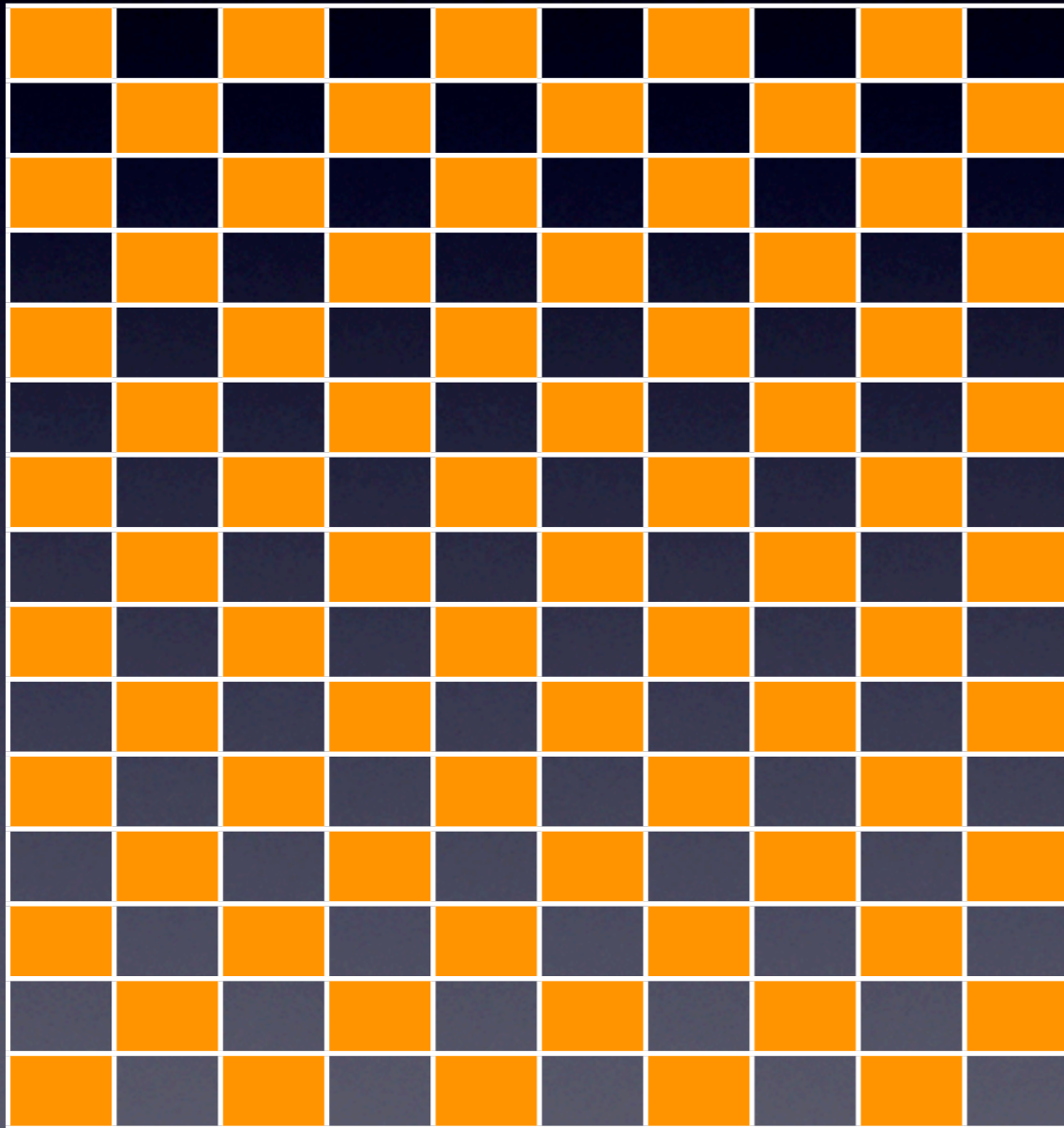
Outline

1. **Memory Management**
2. Concurrency Management
3. I/O Management

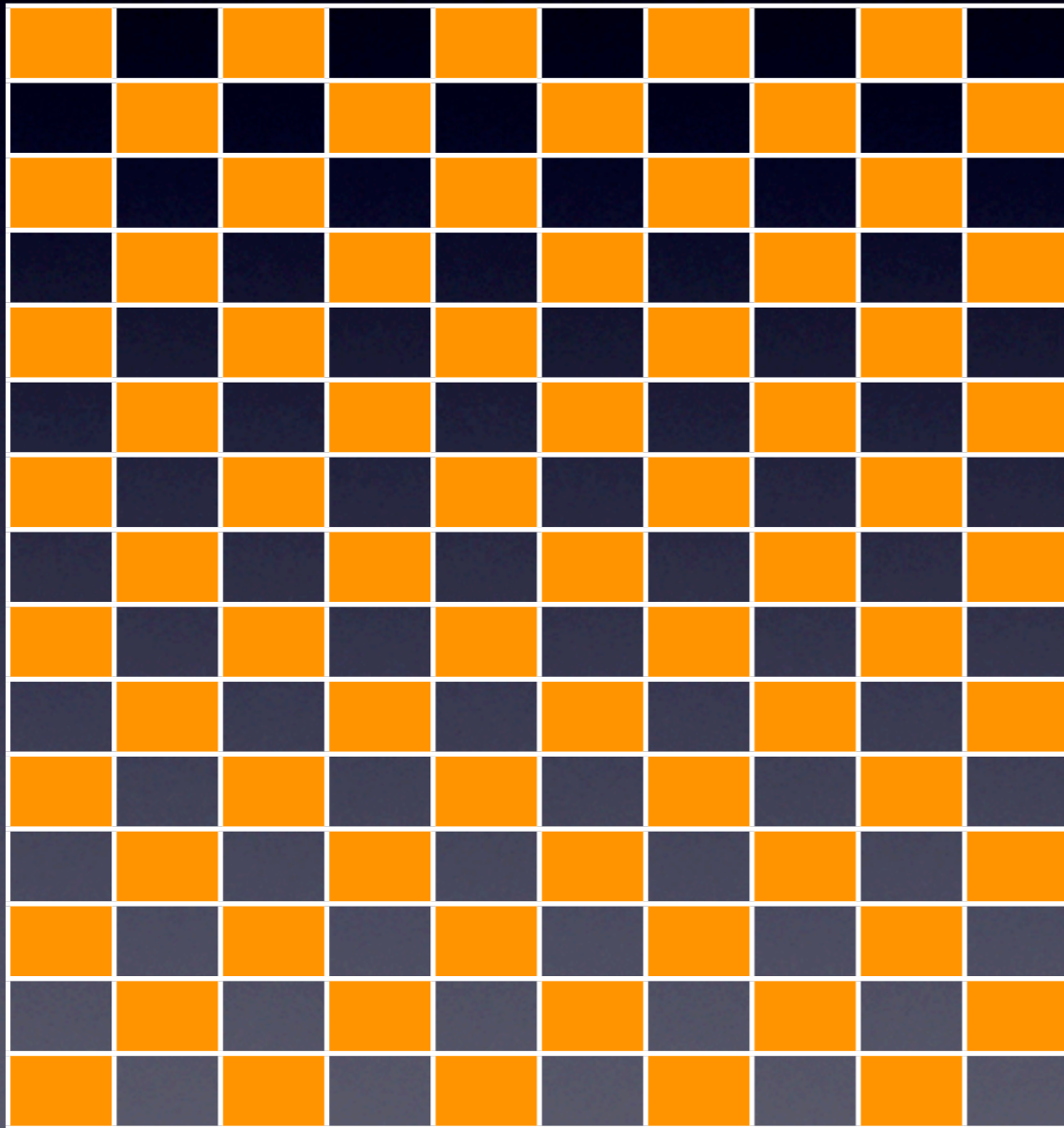
Goals

- `malloc(n)` takes at most $\text{time}(n)$
- `free(n)` takes at most $\text{time}(n)$
- access takes **small** constant time
- **small** and **predictable** memory fragmentation bound

The Problem

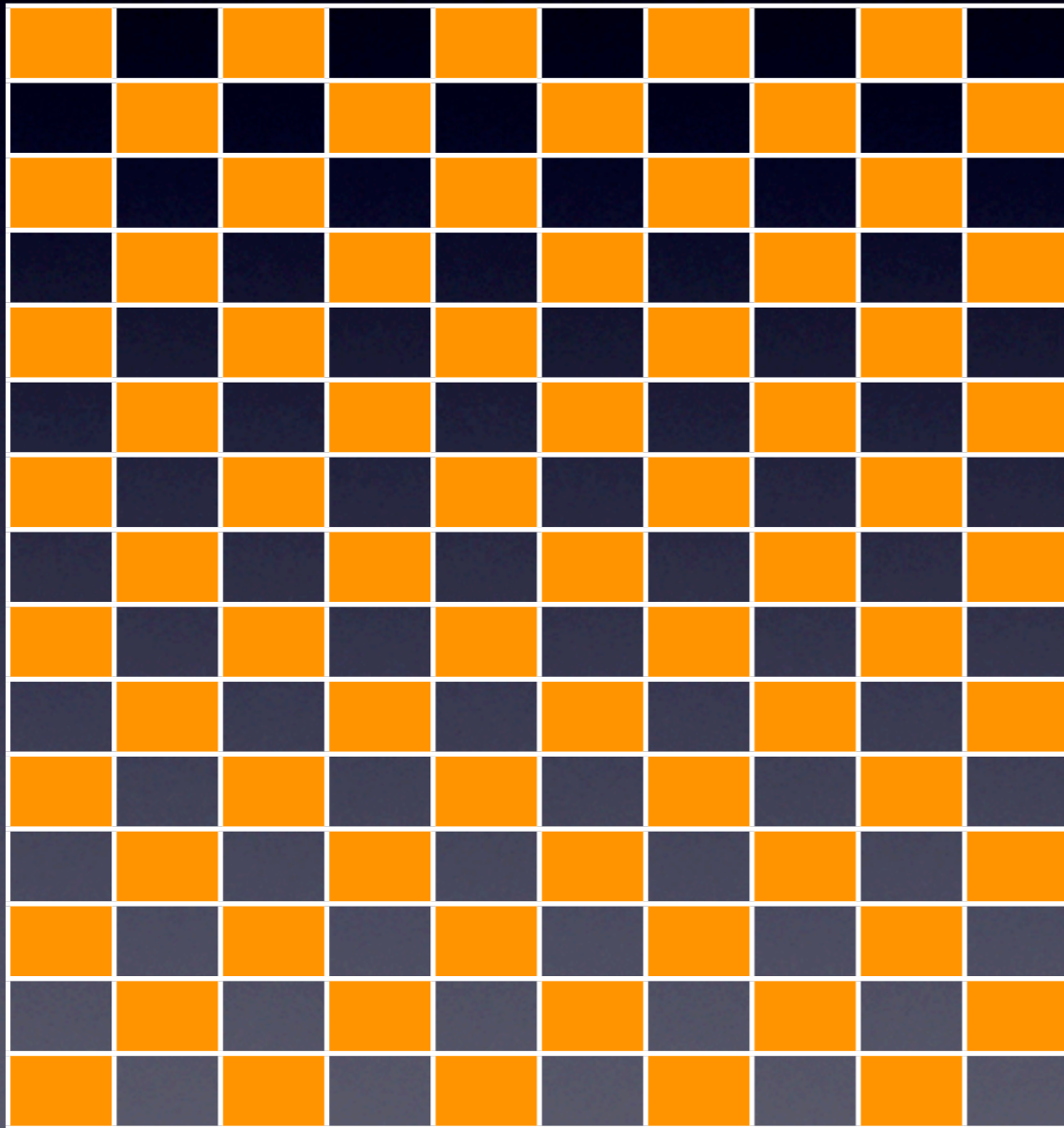


The Problem



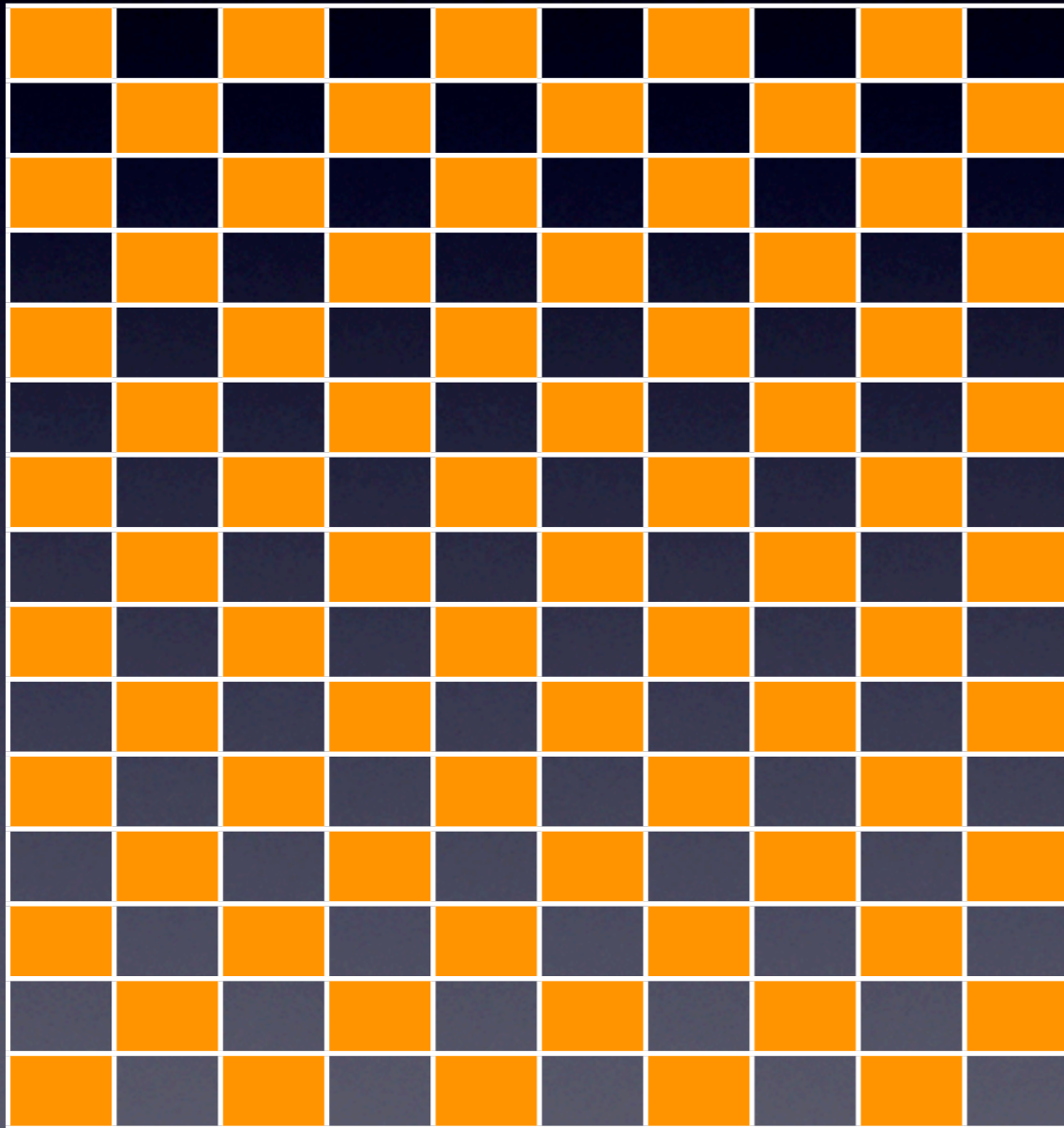
- Fragmentation

The Problem



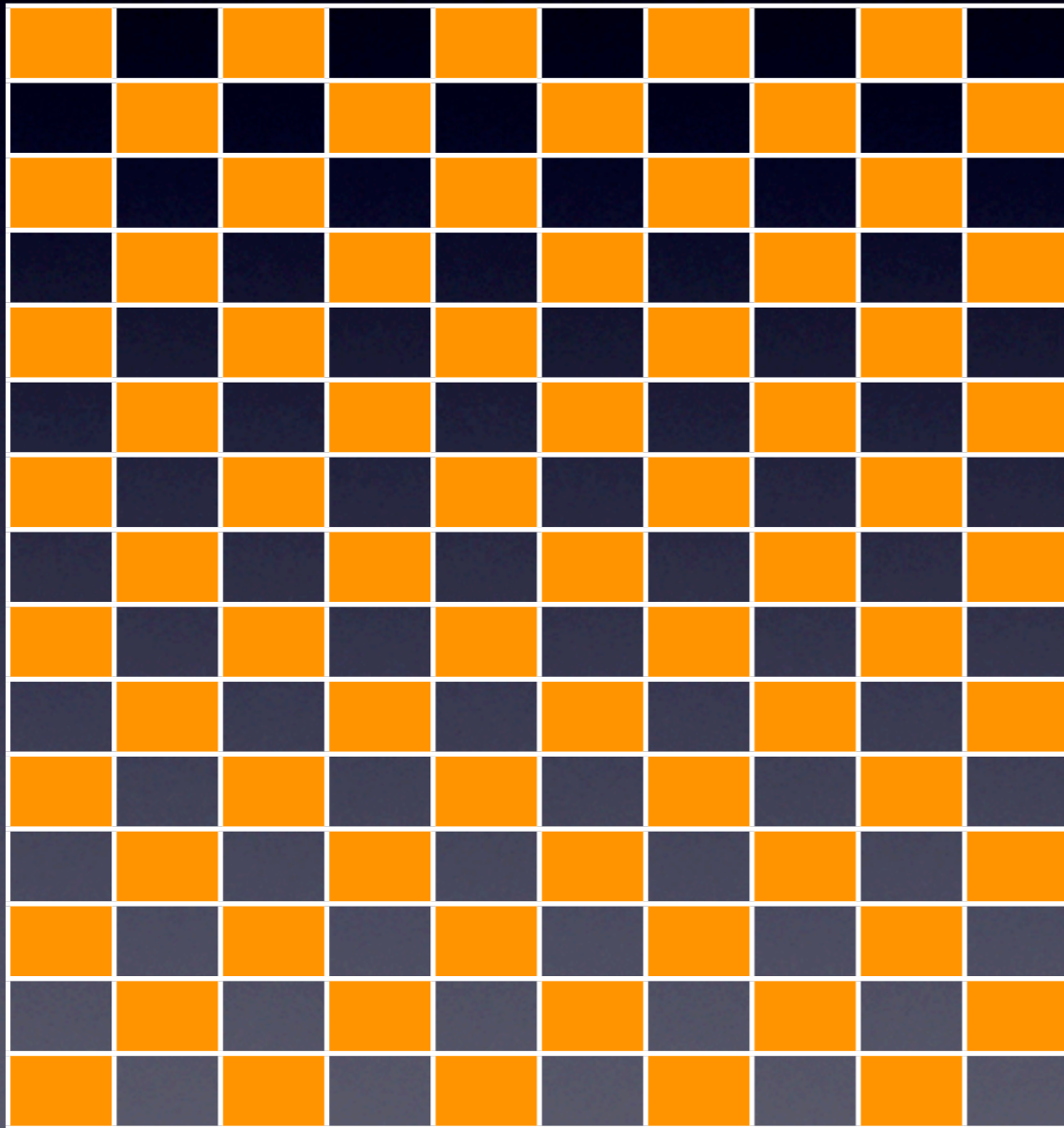
- Fragmentation
 - ▶ Compaction

The Problem

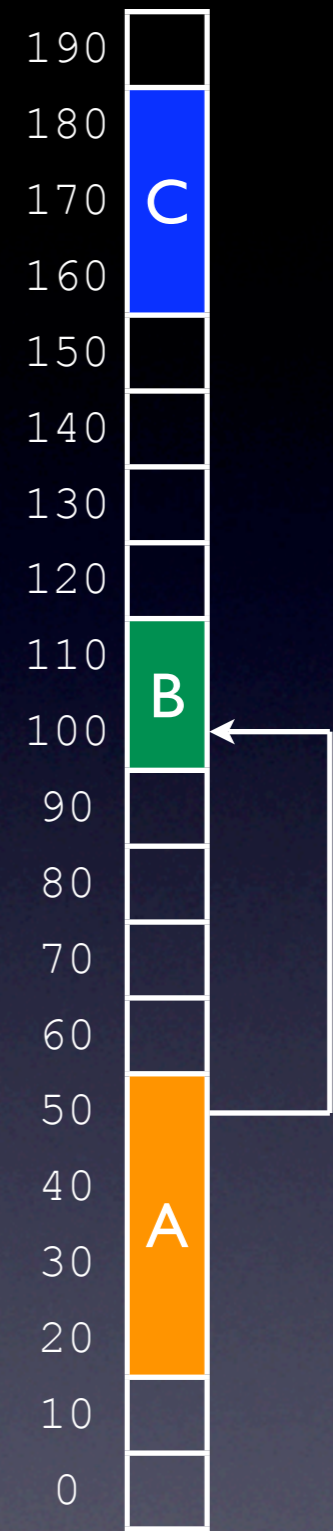


- Fragmentation
 - ▶ Compaction
 - References

The Problem

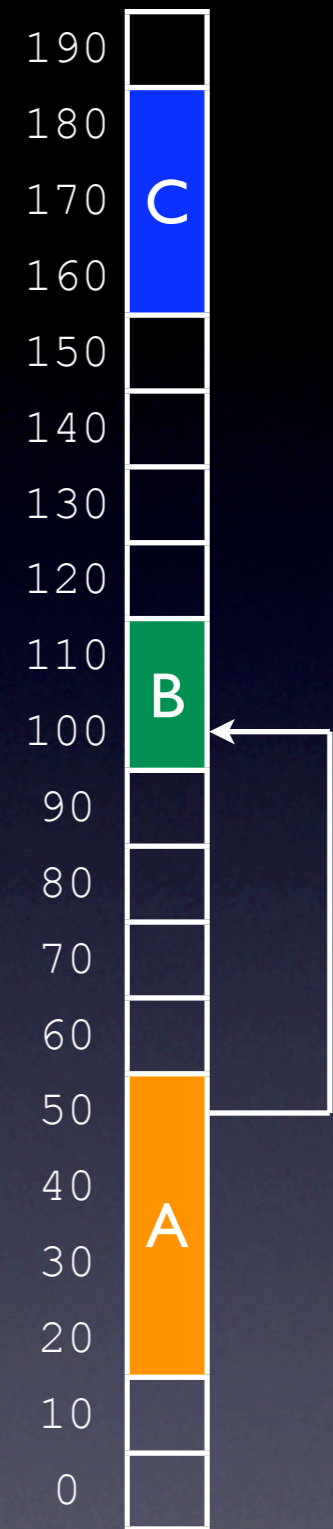


- Fragmentation
 - ▶ Compaction
- References
 - ▶ Abstract Space



Memory

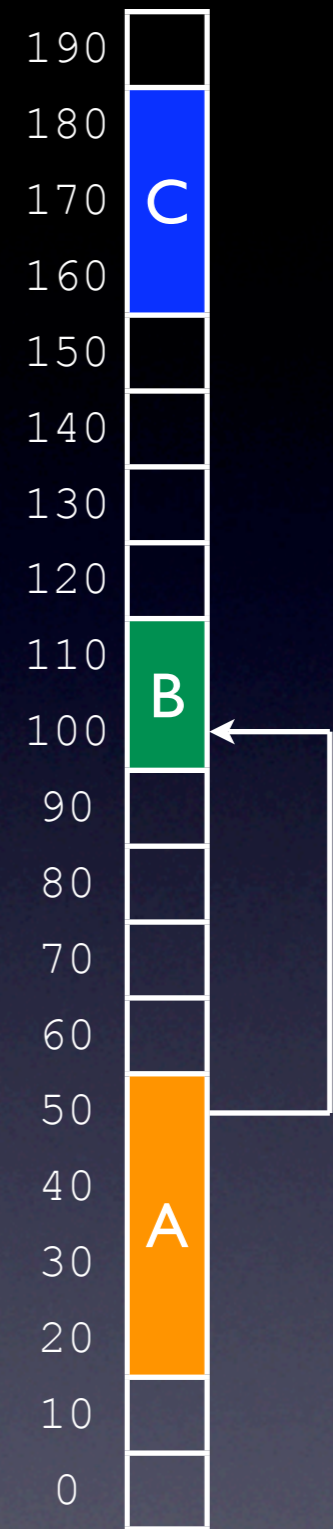
Example:



Memory

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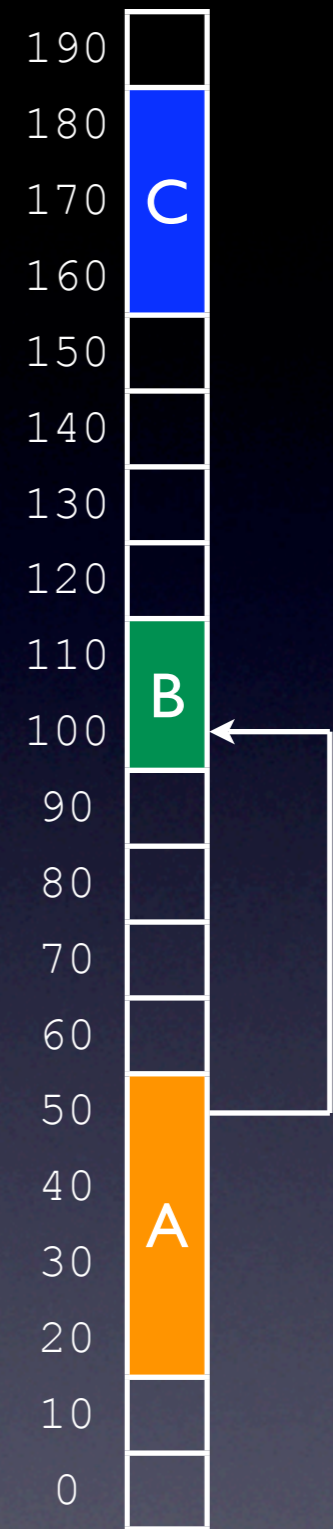
- There are three objects



Memory

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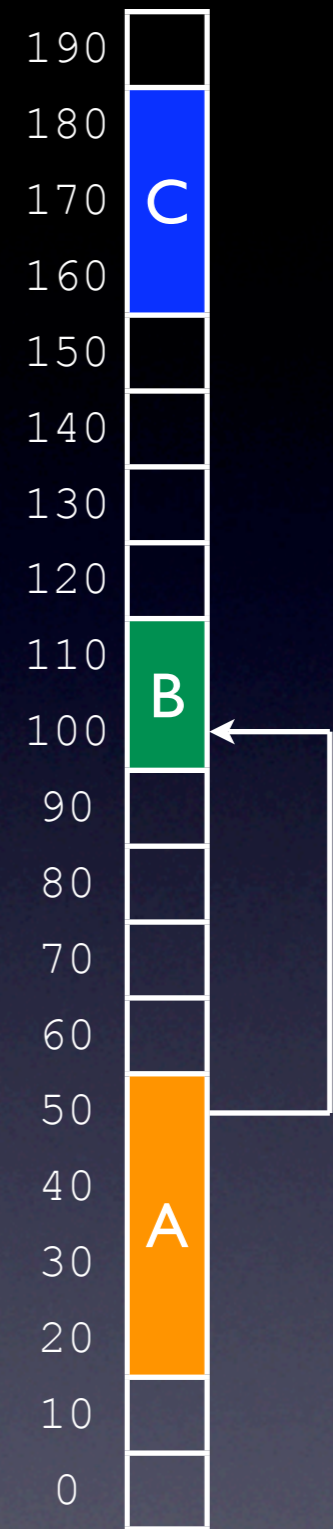
- There are three objects
- Object **A** starts at address 20



Memory

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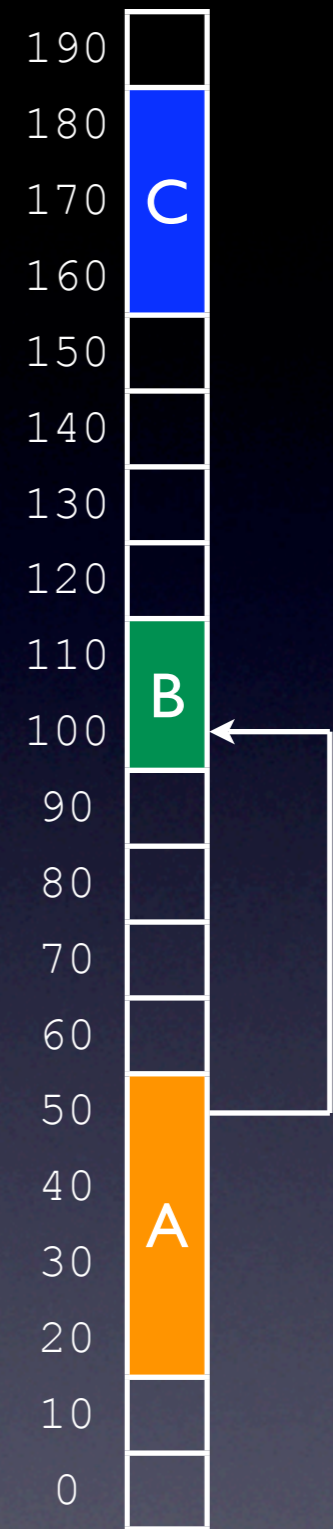
- There are three objects
- Object **A** starts at address 20
- Object **A** needs 40 bytes



Memory

Example:

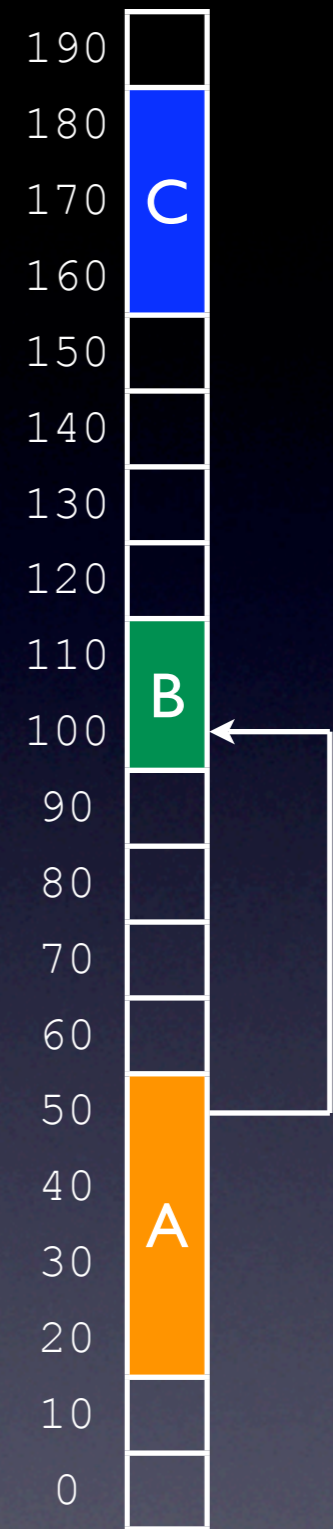
- There are three objects
- Object **A** starts at address 20
- Object **A** needs 40 bytes
- **B** starts at 100, needs 20 bytes



Memory

Example:

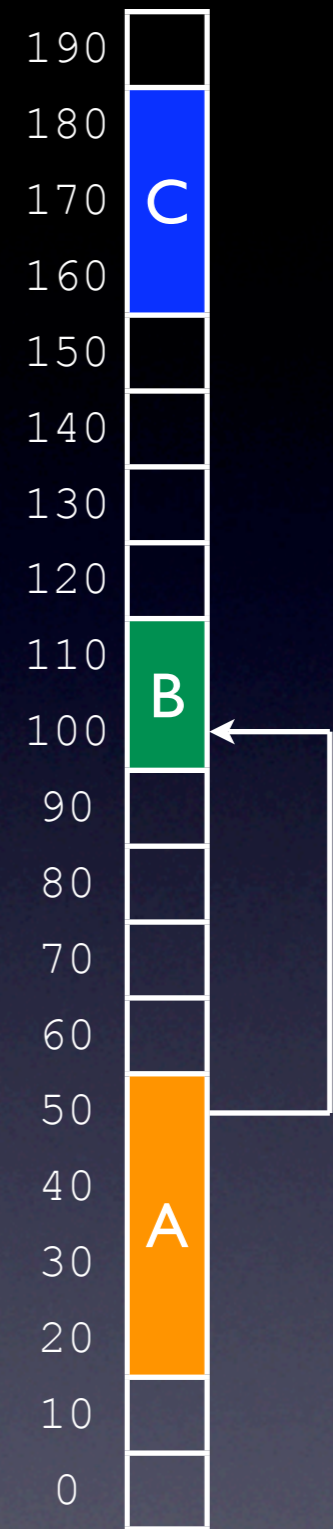
- There are three objects
- Object **A** starts at address 20
- Object **A** needs 40 bytes
- **B** starts at 100, needs 20 bytes
- **C** starts at 160, needs 30 bytes



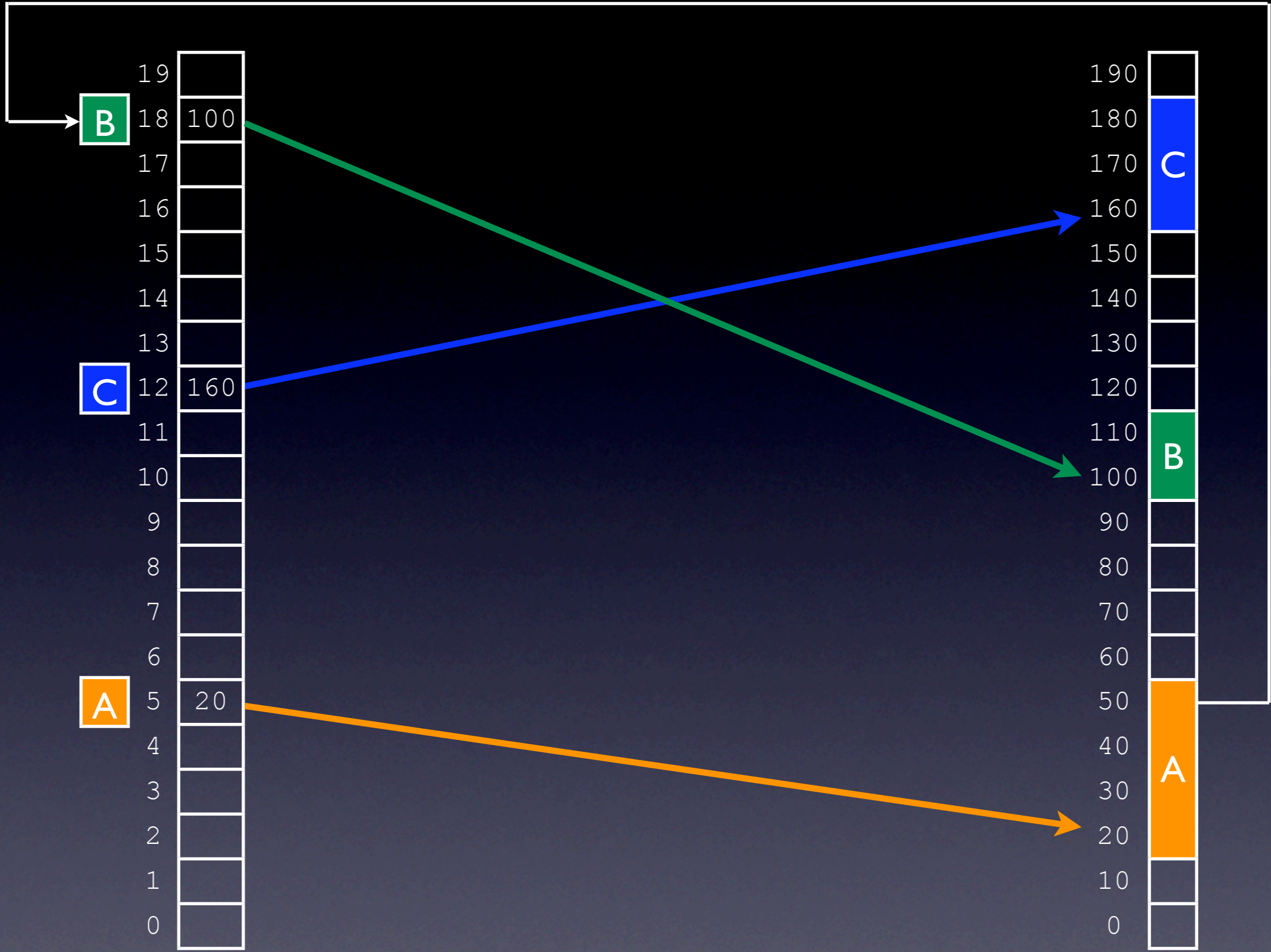
Memory

Example:

- There are three objects
- Object **A** starts at address 20
- Object **A** needs 40 bytes
- **B** starts at 100, needs 20 bytes
- **C** starts at 160, needs 30 bytes
- **A** contains a reference to **B**



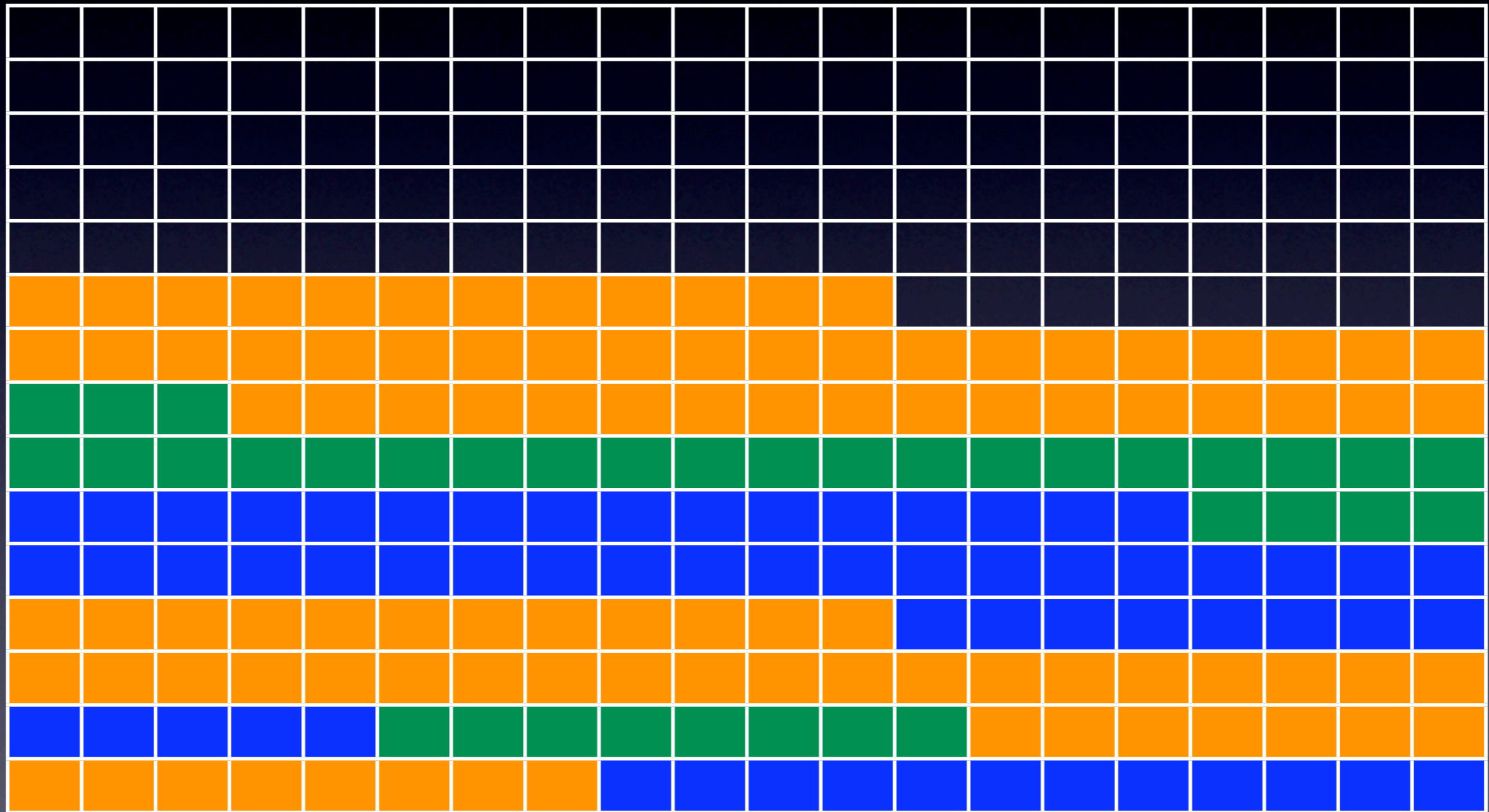
Memory



Abstract Space

Concrete Space

Keep It Compact?



Trade-Off Speed for Memory Fragmentation

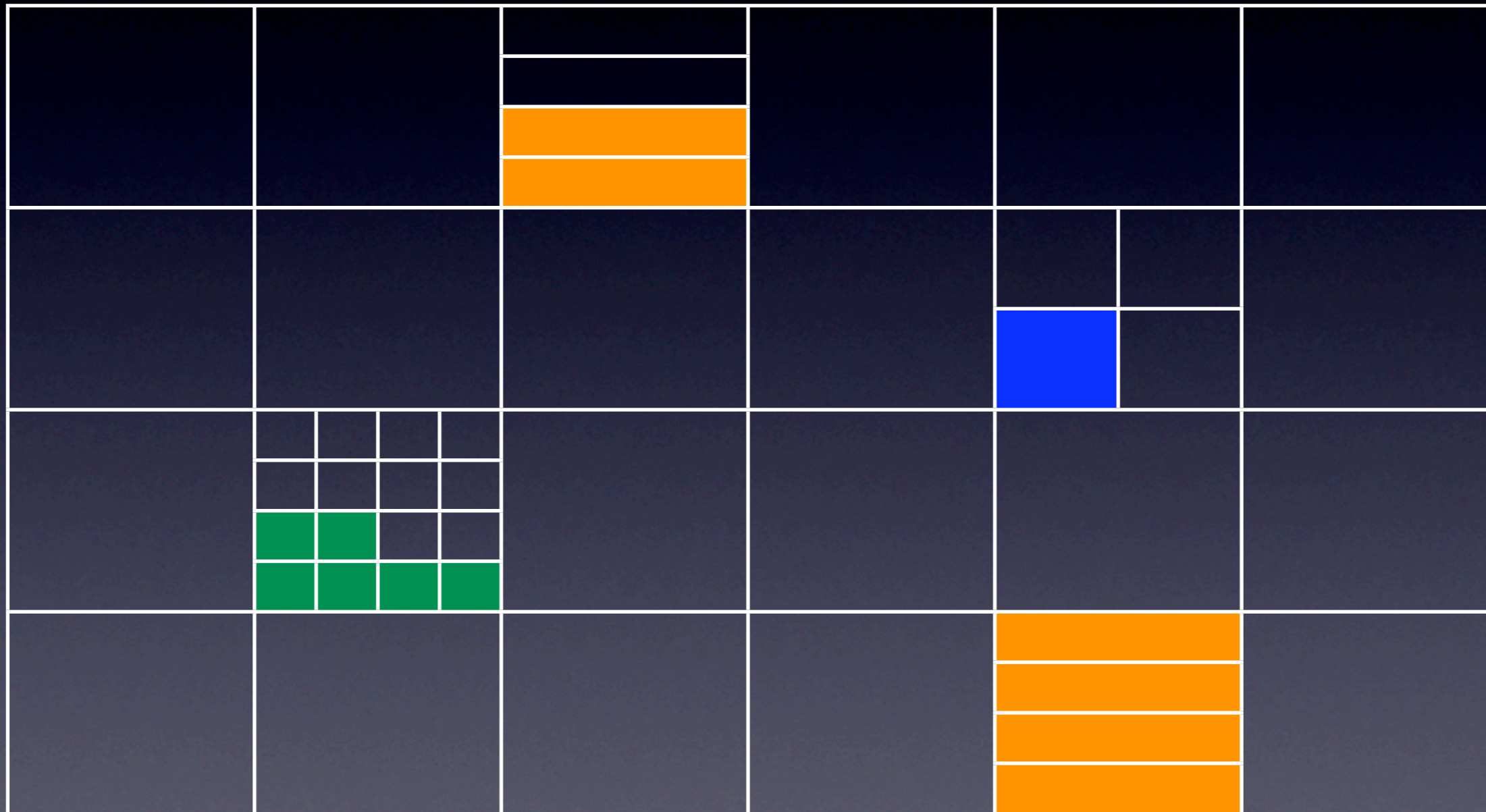
Trade-Off Speed for Memory Fragmentation

Keep Speed and
Memory Fragmentation
Bounded and **Predictable**

Partition Memory into Pages

16KB	16KB	16KB	16KB	16KB	16KB
16KB	16KB	16KB	16KB	16KB	16KB
16KB	16KB	16KB	16KB	16KB	16KB
16KB	16KB	16KB	16KB	16KB	16KB

Partition Pages into Blocks



Size-Class Compact



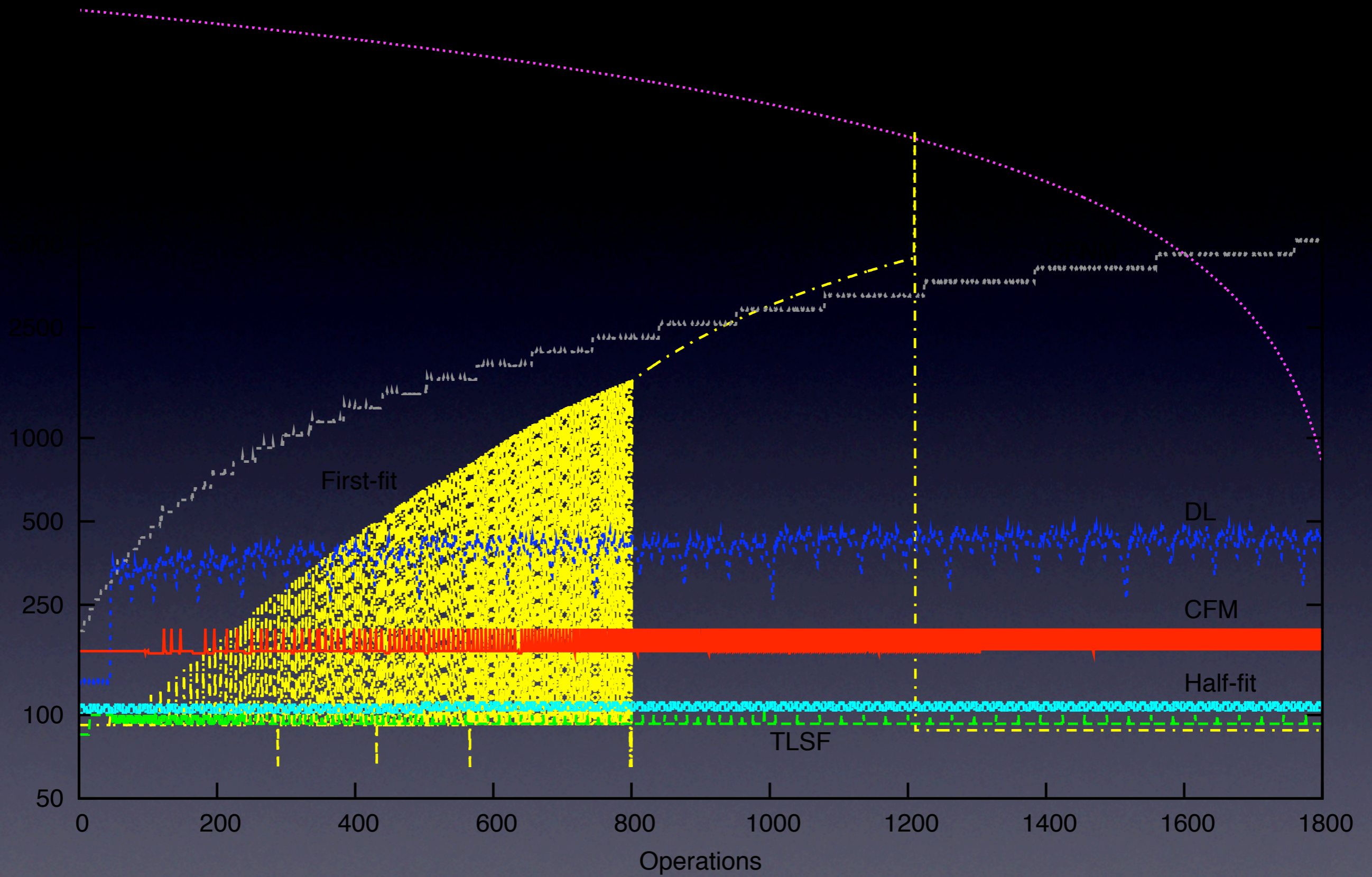
Objects < 32



Objects < 48



Objects < 64

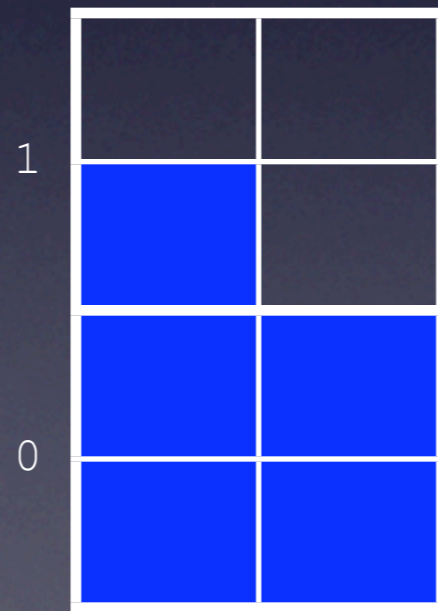


Bounded Compaction

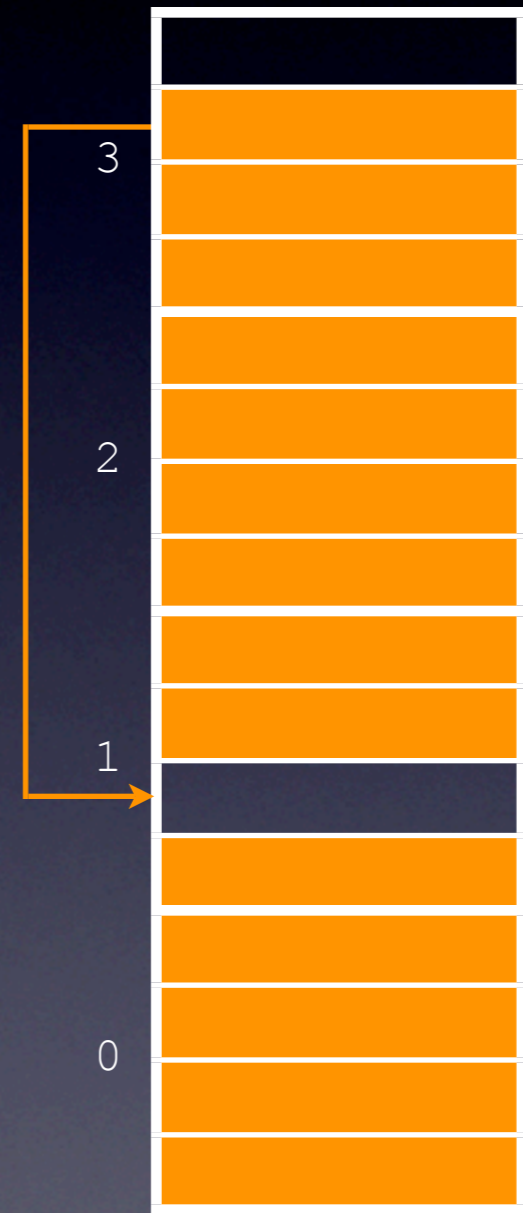
just **move** 'last' object



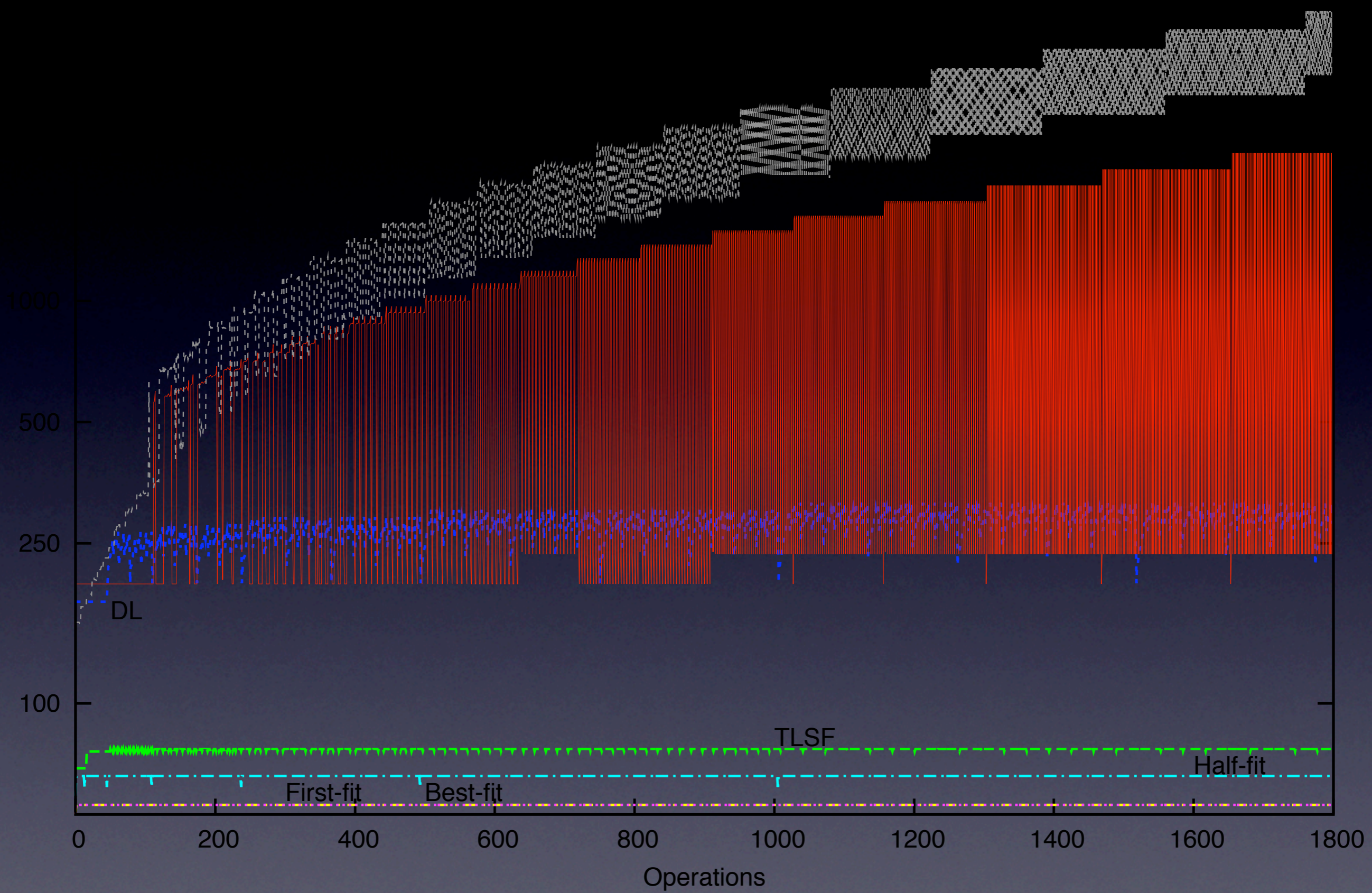
Objects < 32



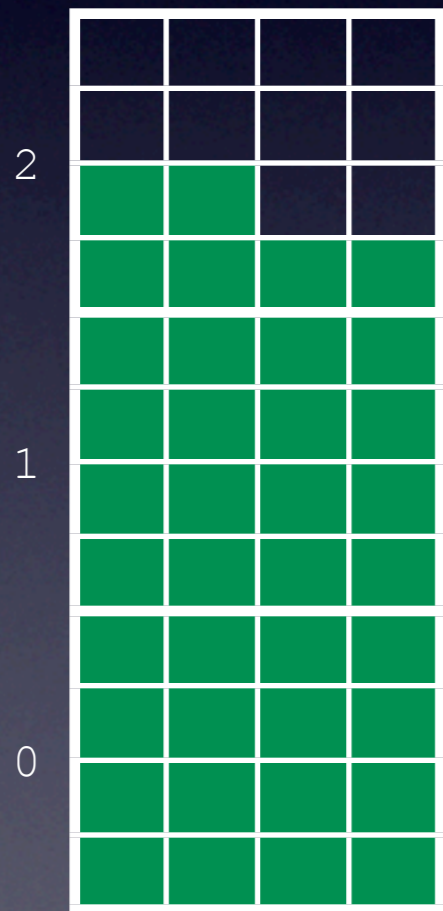
Objects < 48



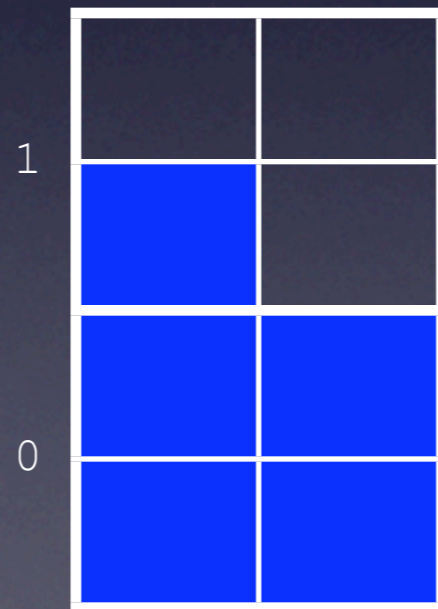
Objects < 64



Partial Compaction



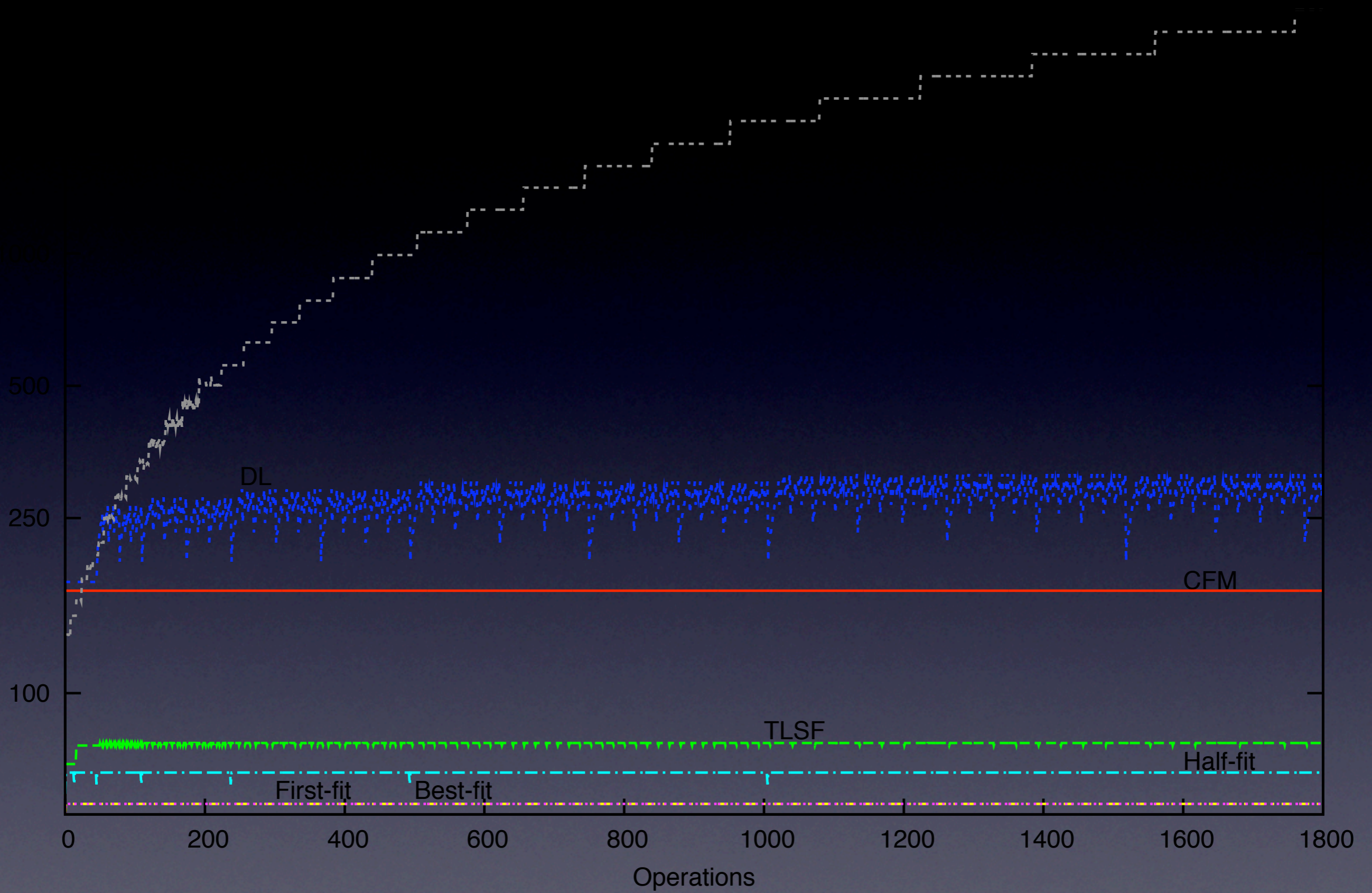
Objects < 32



Objects < 48



Objects < 64



Program Analysis

Program Analysis

Definition:

Let k count deallocations in a given size-class for which no subsequent allocation was done (“k-band mutator”).

Program Analysis

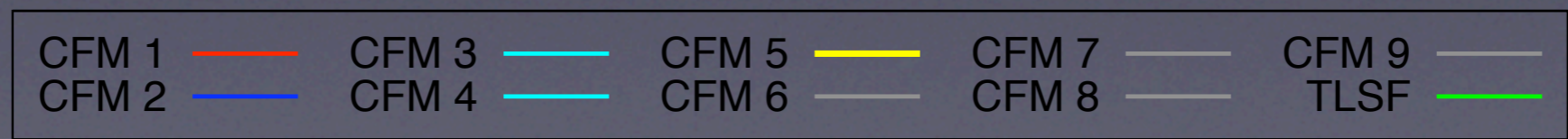
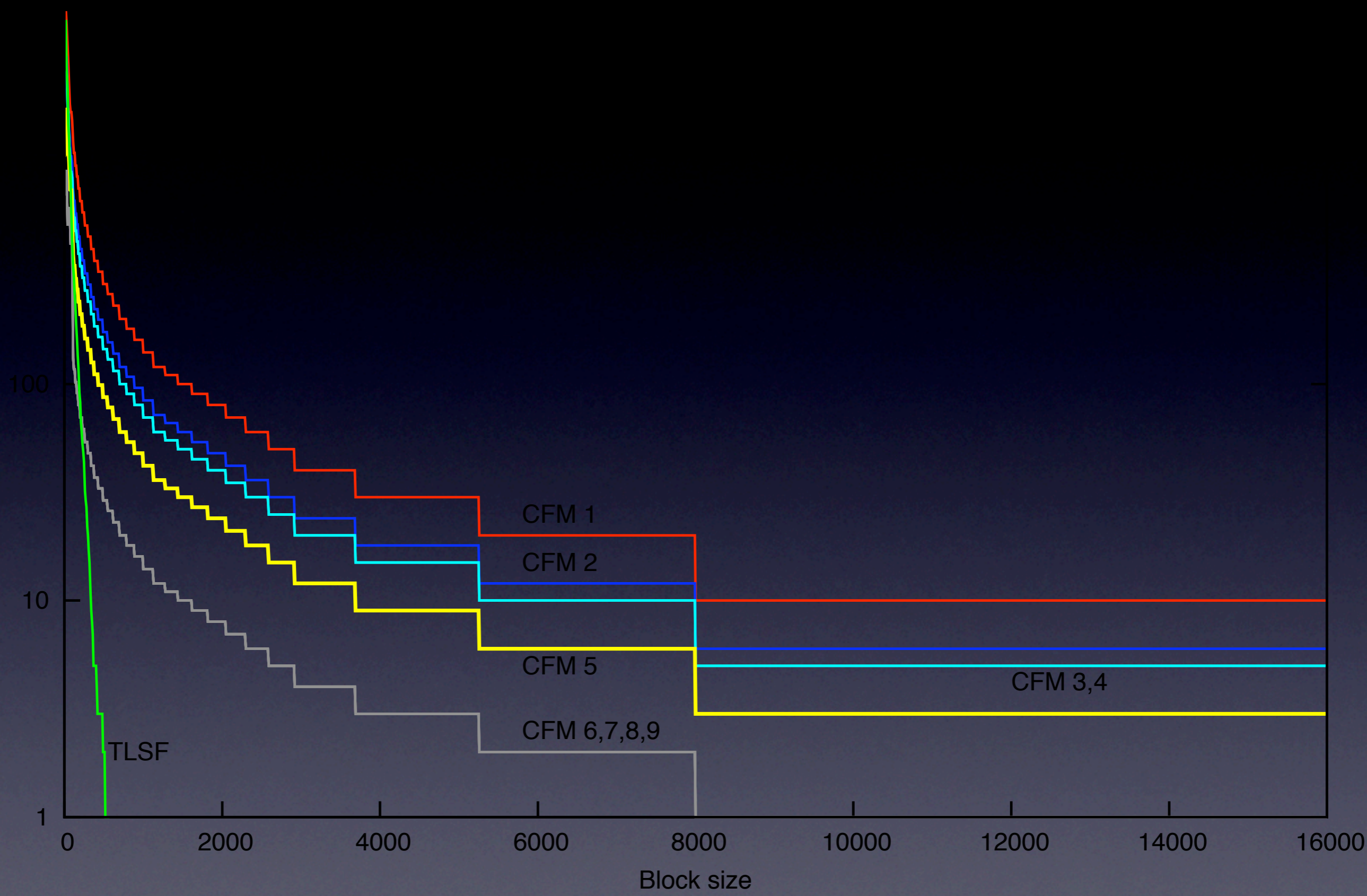
Definition:

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Proposition:

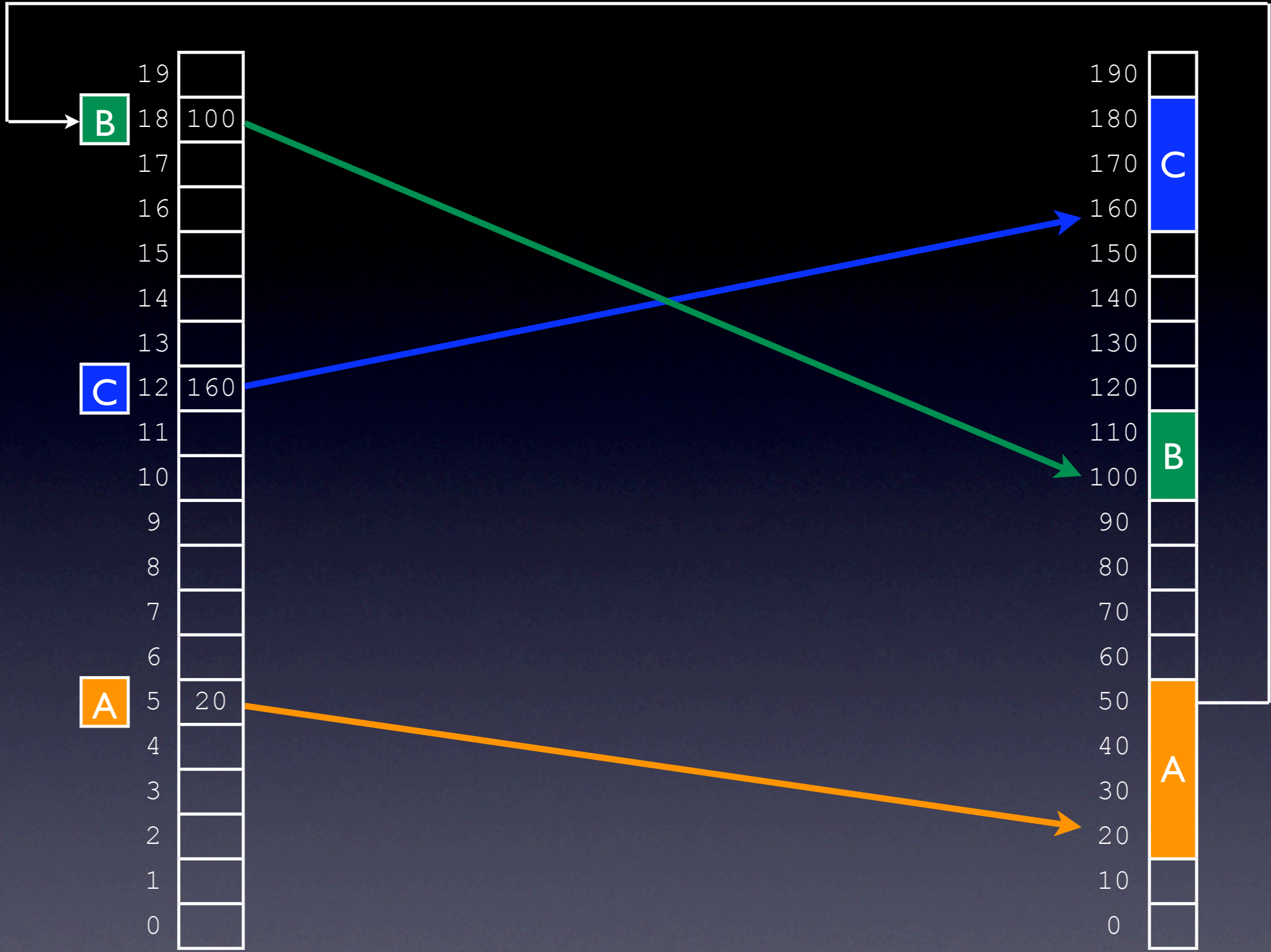
Each deallocation that happens when $k < \text{max_number_of_non_full_pages}$ takes constant time.

Number of allocatable blocks



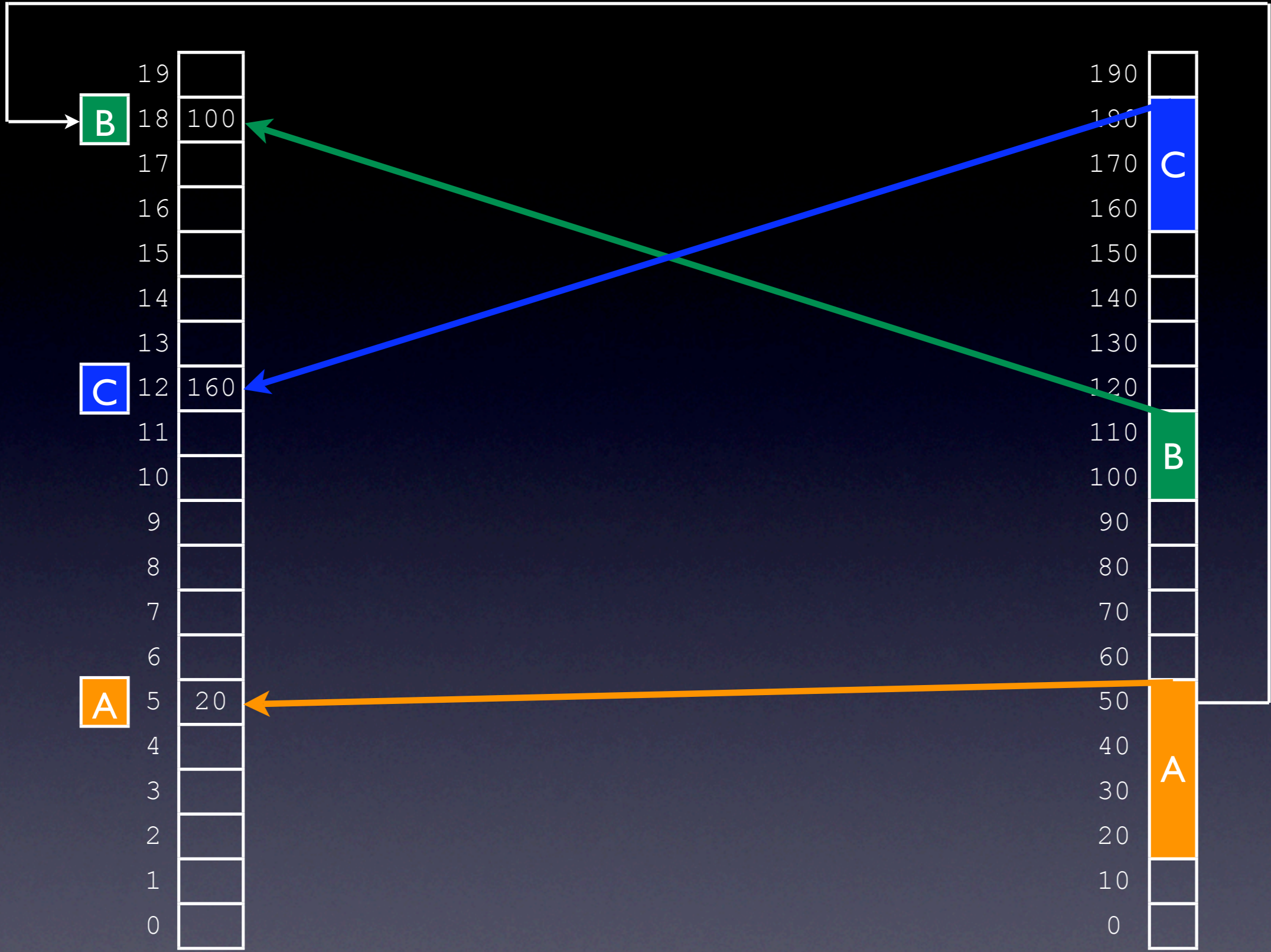
Results

- if mutator stays within k-bands:
 - malloc(n) takes constant time
 - free(n) takes constant time
 - access takes one indirection
- memory fragmentation **bounded** in k and **predictable** in constant time



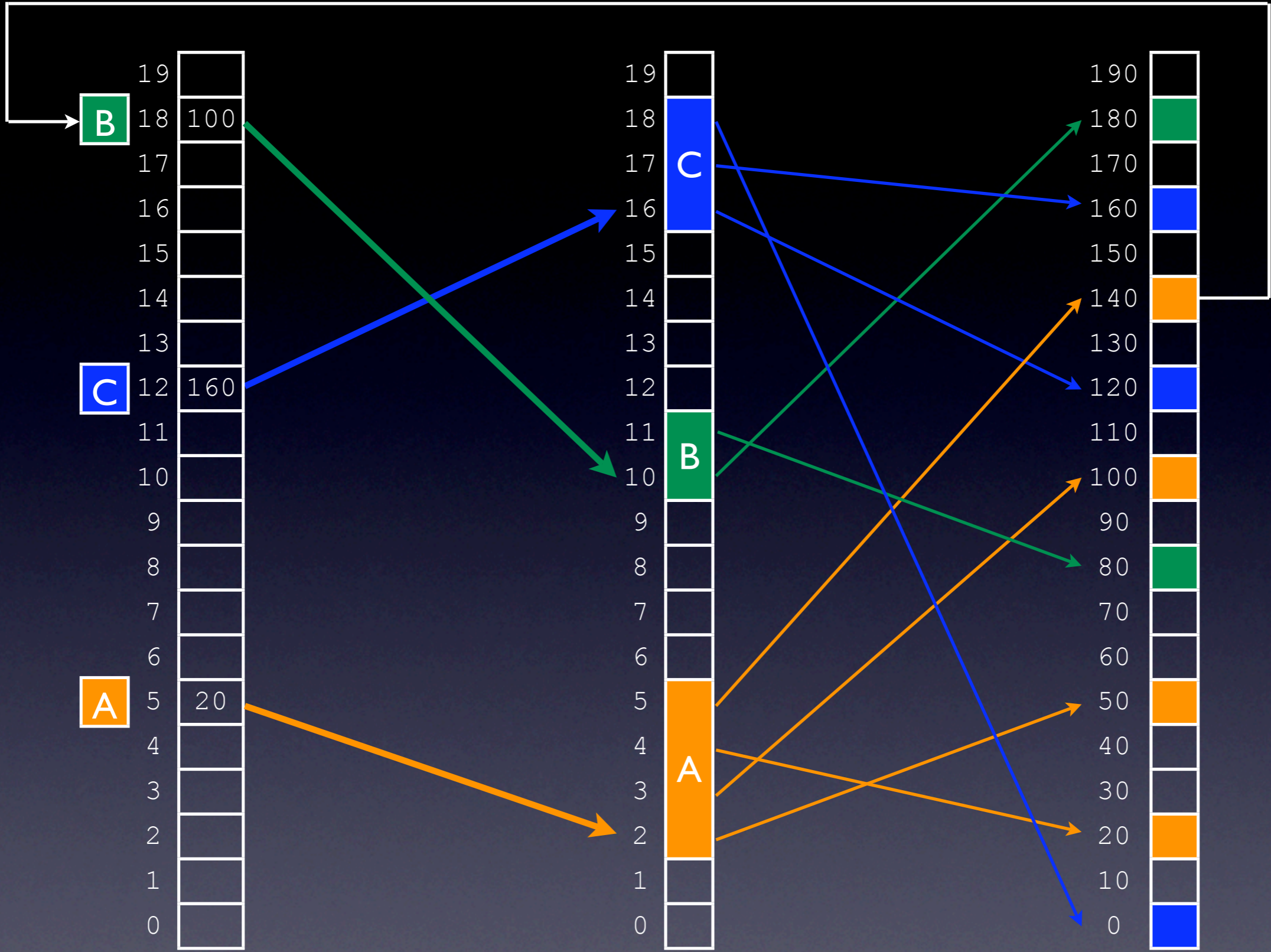
Abstract Space

Physical Memory



Abstract Space

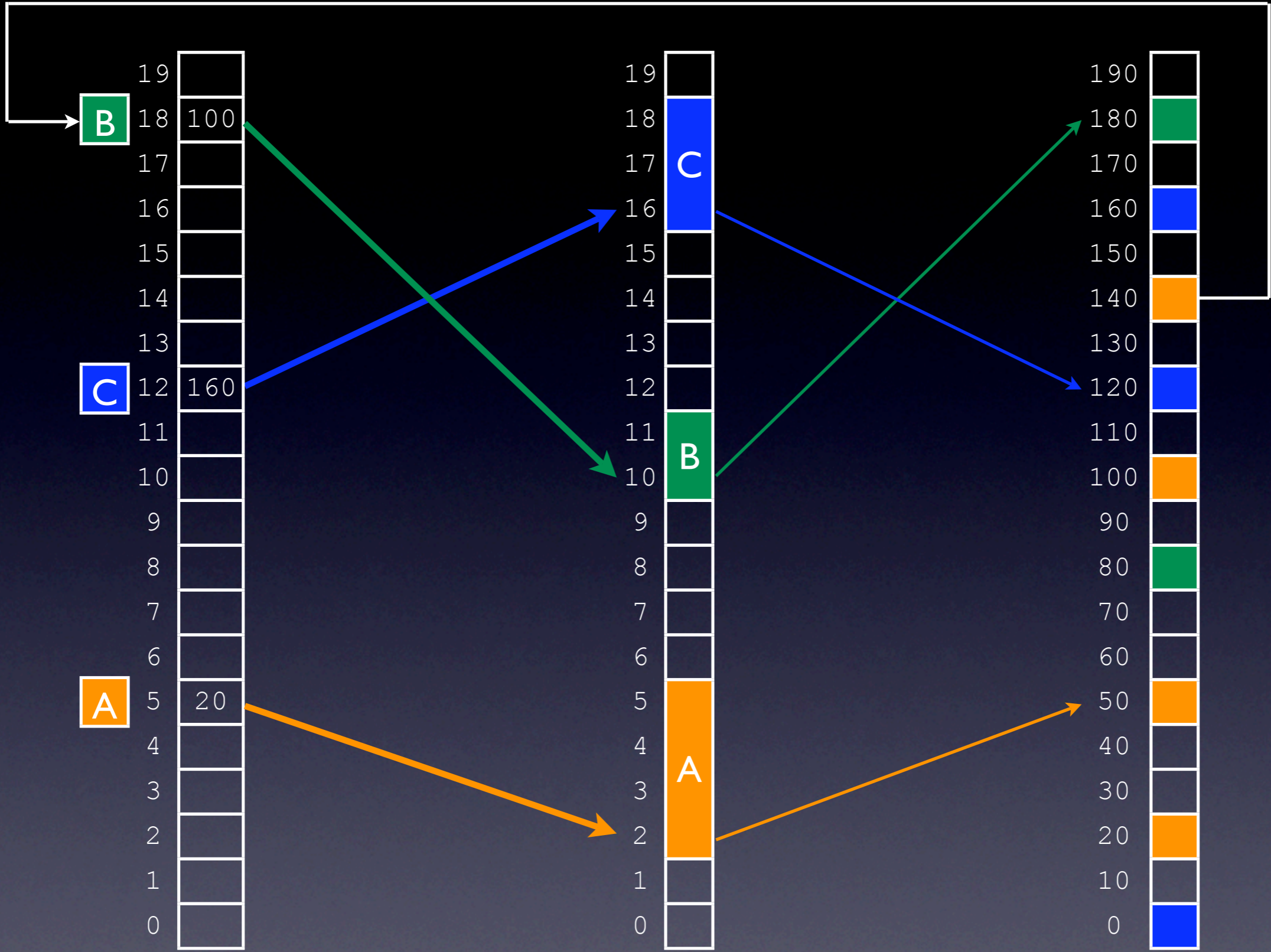
Physical Memory



Abstract Space

Virtual Space

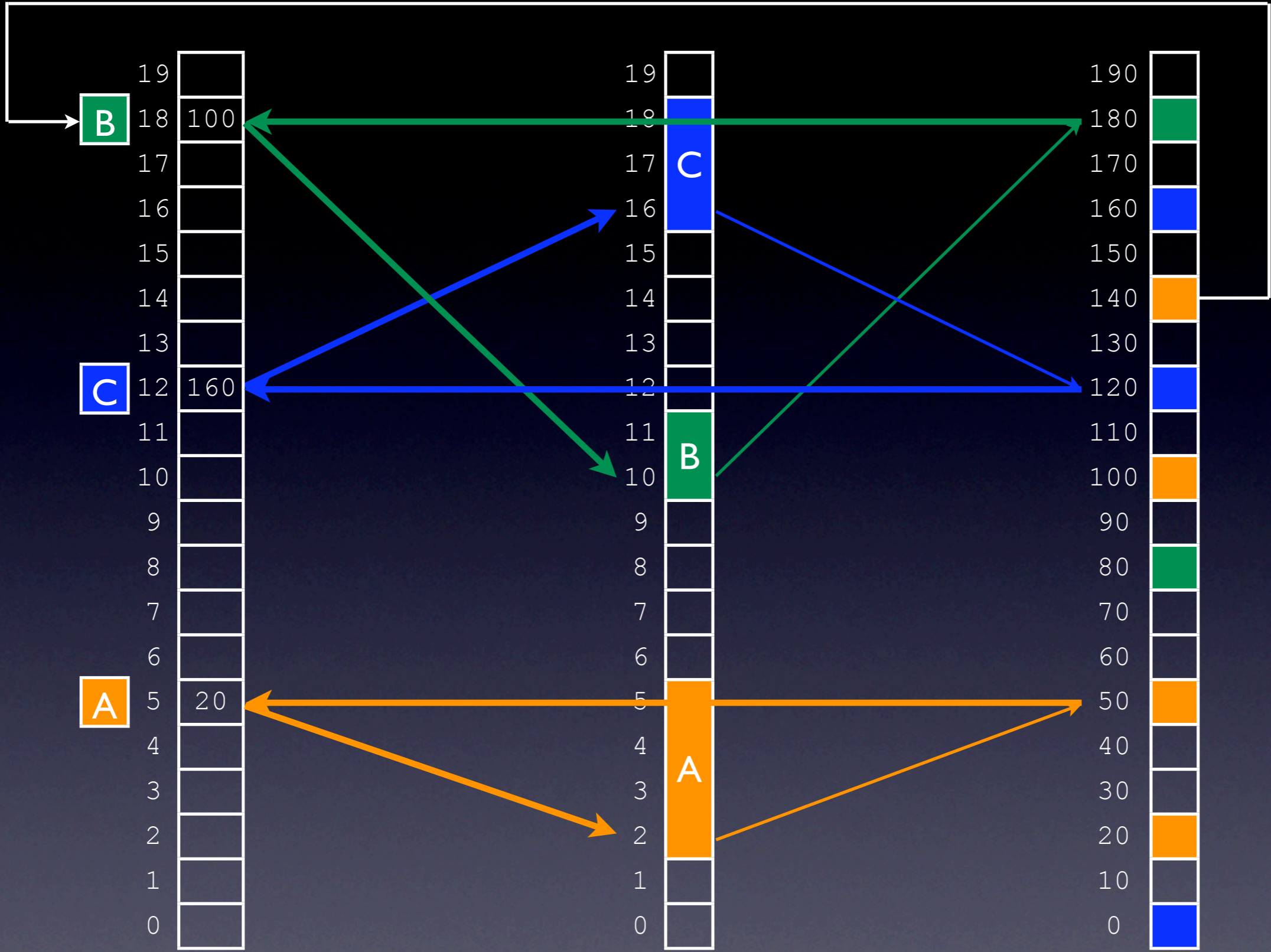
Physical Memory



Abstract Space

Virtual Space

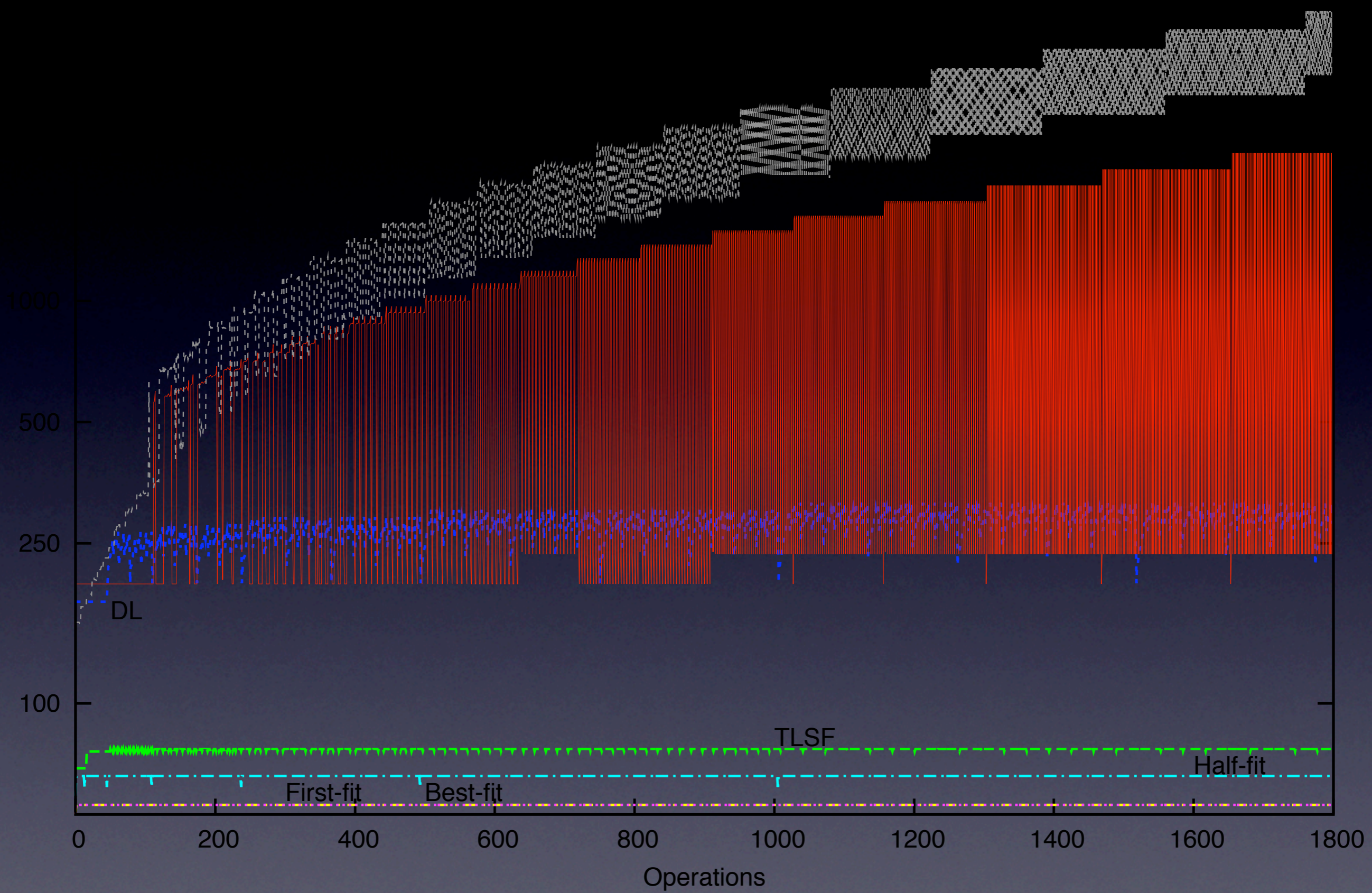
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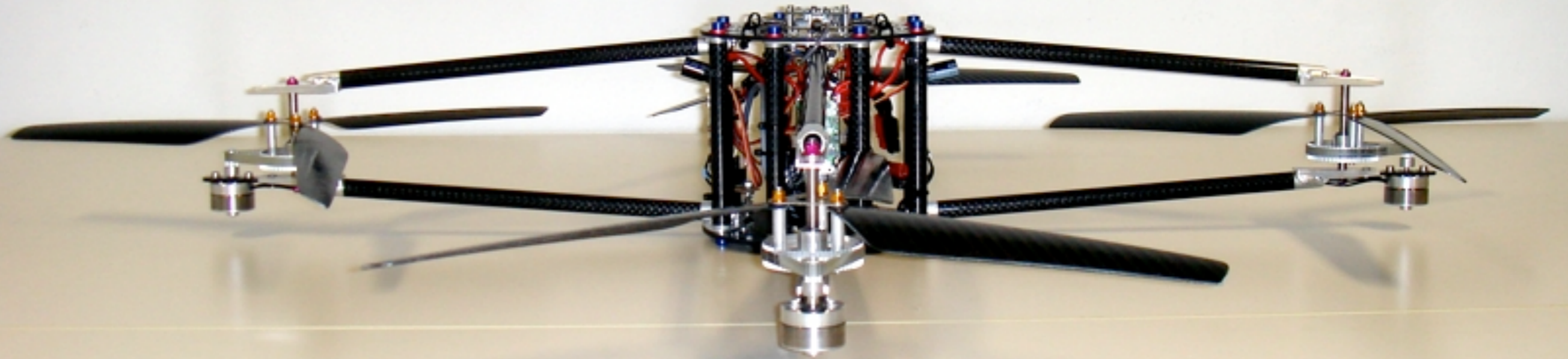


Abstract Space

Virtual Space

Physical Memory





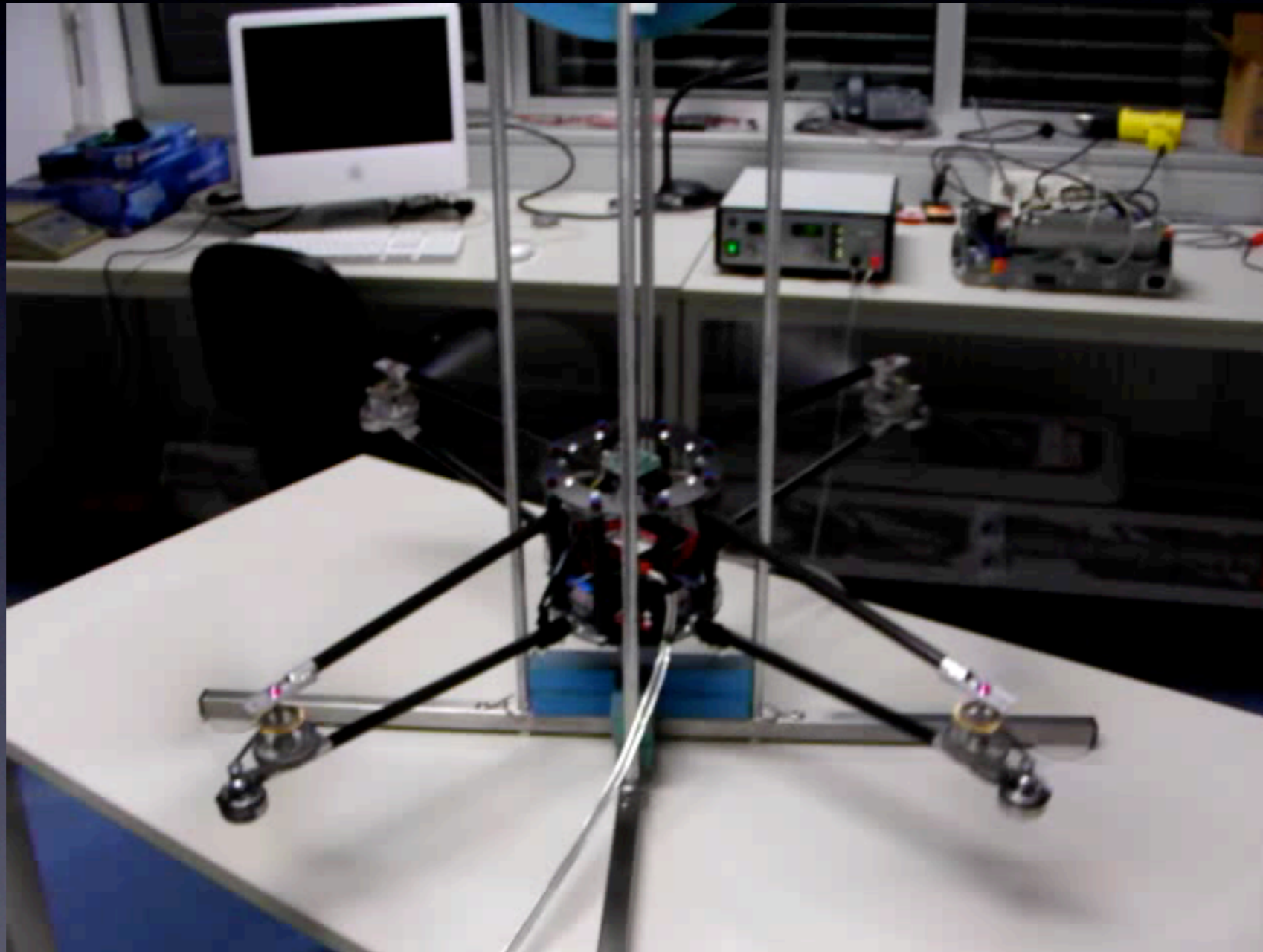
The JAviator

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Quad-Rotor Helicopter



Flight Control



Thank you