# The Logical Execution Time Paradigm

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#### ESWEEK Tutorial, Taipei, October 2011

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## Control Software



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A LET Program Incorporates Logical Execution Time as First-Class Concept

### Definition

### Plant

A LET program's I/O behavior is input-determined if, for all sequences I of input values and times, the system always produces unique sequences f(I) of output values and times.



### Result

### Plant

A LET program's I/O behavior is input-determined on any platform that runs the program time-safely.

[with Henzinger, Horowitz at EMSOFT 2001 and in Proceedings of the IEEE 2003]



# Giotto on ETHZ Helicopter, 2001

# From Control Models to Real-Time Code Using Giotto

[with Henzinger, Sanvido, Pree in the IEEE Control Systems Magazine, 2003]

## HTL:

A Hierarchical Coordination Language for Interacting Real-Time Tasks

[with Ghosal, Henzinger, Iercan, and Sangiovanni-Vincentelli at EMSOFT 2006]

# Design: Adding Tasks

Plant



# Design: Adding Tasks



### Observation

Plant

A LET program's existing I/O behavior does not change by adding new tasks.

# Analysis: Refining Tasks



# Analysis: Refining Tasks



### Result

Plant

A concrete LET program is time-safe if it refines a time-safe, abstract LET program.

[with Ghosal, Henzinger, Iercan, Sangiovanni-Vincentelli at EMSOFT 2006]

## Distributed, Modular HTL

[with Henzinger, Marques, Sokolova at RTSS 2009]

# Complexity

$\varphi$	C	$\mathcal{D}^{\mathcal{A}}_{\varphi}(C,P)$	$\mathcal{C}^{\mathcal{A}}_{\varphi}(C,P)$	$\overline{\mathcal{C}}_{arphi}^{\mathcal{A}}(P)$
Well-formedness	any	C	$n_{m\downarrow}^C \; n_T \; n_p$	$n_{m\downarrow}^P \; n_T \; n_p$
Race freedom	top	P	$n_{T\uparrow}^C n_w + n_M n_c$	$n_{T\uparrow}^P n_w + n_M n_c$
	ref.	C	1	
Transmission safety	any	C	1	$n_c$
Time safety	top	P	$(n_m \ \Delta_{max})^{n_M}$	$(n_m \Delta_{max})^{n_M}$
	ref.	C	1	
Code generation	any	C	$n_{m\downarrow}^C \left( n_T  n_a + n_m \right)$	$n_{m\downarrow}^P \left( n_T  n_a + n_m \right)$

$n_a$	number of communicator accesses per task
$n_m$	number of modes per module
$n_w$	number of communicator writes per task

 $n_{T\uparrow}^C$ number of top-level tasks in C

number of communicators  $n_c$  $n_p \ n_{m\downarrow}^C \ n_{T\uparrow}^P$ number of ports per task

- total number of modes in C
- number of top-level tasks in P

number of modules per program  $n_M$  $n_T \\ n_{m\downarrow}^P$ number of tasks per mode total number of modes in Pmaximal value of mode periods  $\Delta_{\max}$ 

### Result I

Plant

#### Checking time-safety (and transmission-safety) of HTL programs is <u>modular</u> (below top level).

[with Henzinger, Marques, Sokolova at RTSS 2009]

## Result II

Plant

## Generating distributed code for HTL programs is <u>modular</u>.

[with Henzinger, Marques, Sokolova at RTSS 2009]



### The JAviator javiator.cs.uni-salzburg.at

# Quad-Rotor Helicopter



• all carbon, titanium, aluminum design

• custom motors

I.3m diameter
~2.2kg weight
+2kg payload

~40min (empty)
~10min (full)

#### [AIAA GNC 2008]

# Open Source Blueprints





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# Outdoor Flight Salzburg Controller



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# Thank you

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