

Scalability of Vehicle Networks through Vehicle Virtualization

Jiangchuan Huang*, Christoph Kirsch† and Raja Sengupta* *Systems Engineering, Department of Civil and Environmental Engineering, University of California, Berkeley. †Department of Computer Sciences, University of Salzburg.

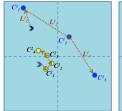
Lab website: cpcc.berkeley.edu. Email: jiangchuan@berkeley.edu. Presented by Elói Pereira*

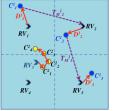
Cyber-Physical Applications:

- Google Street View.
- Real-time traffic reporting.
- Unmanned aerial vehicle (UAV)-based sensing.
- Mobile sensor networks.

Virtual Vehicle (VV):

- A replica of a real vehicle (RV).
 - 1. Physical-mobility: travels (binding) with an RV.
 - 2. Cyber-mobility: migrates from one RV to another.



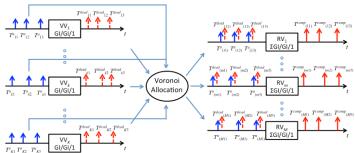


(a) Customer view

(b) Provider view

- Theorem 1: Each VV is a GI/GI/1 queuë.
- Provider aims to complete each task before virtual departure time.
- The virtual departure time is like a deadline for each task, thus called virtual deadline.
- Virtual vehicle and virtual deadline create a soft real-time system.
 - Tardiness = max{Actual completion time Virtual deadline, 0}
 - Delivery prob = Pr (Actual completion time ≤ Virtual deadline)
- High performance isolation if a statistically dominant subset [1], e.g., Delivery prob = 98%, of the virtual vehicle's tasks are completed no later than their virtual deadline.

Voronoi Allocation: Theorem 2: Each RV is a ΣGI/GI/1 queue.



Gain: # VVs / # RVs.

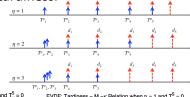
- Multiplexing gain: a customer may not fully utilize her VV.
- Migration gain: gain by migrating the VV to an RV closer to the task.

Scheduling: Earliest Virtual Deadline First (EVDF)

Theorem 3: EVDF achieves minimum tardiness.

Worst-Case Arrival: "n = 1" process.

• Theorem 5: Assume EVDF, the "η = 1" process achieves maximum tardiness among all the renewal processes.



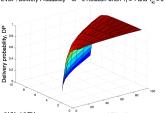
Conclusion:

 High migration gain and high performance isolation.

· Economy of scale: migration gain increases with # RVs.

 Migration gain happens when traveling, not standing still.

 Migration cost is bounded.



VVs / # RVs, ĸ

Migration Cost:

- MC = (size)*(inter distance)*(1{migration}) / (inter virtual departure time).
- MC has unit (bit-meters/second).

[1] Gerald J. Popek and Robert P. Goldberg. Formal requirements for virtualizable third generation architectures. Commun. ACM, 17(7): 412-421, July 1974.

